



Performance Visualizations

Brendan Gregg
Software Engineer

brendan.gregg@joyent.com

G'Day, I'm Brendan

... also known as "shouting guy"

癮科技小劇場：硬碟也會鬧情緒

http://chinese.engadget.com/2009/01/05/video-shouting-at-disk-drive-causes-high-latency-low-r

別向硬碟亂叫，會影響它們的心情...

癮科技小劇場：硬碟也會鬧情緒

癮科技小劇場：硬碟也會鬧情緒

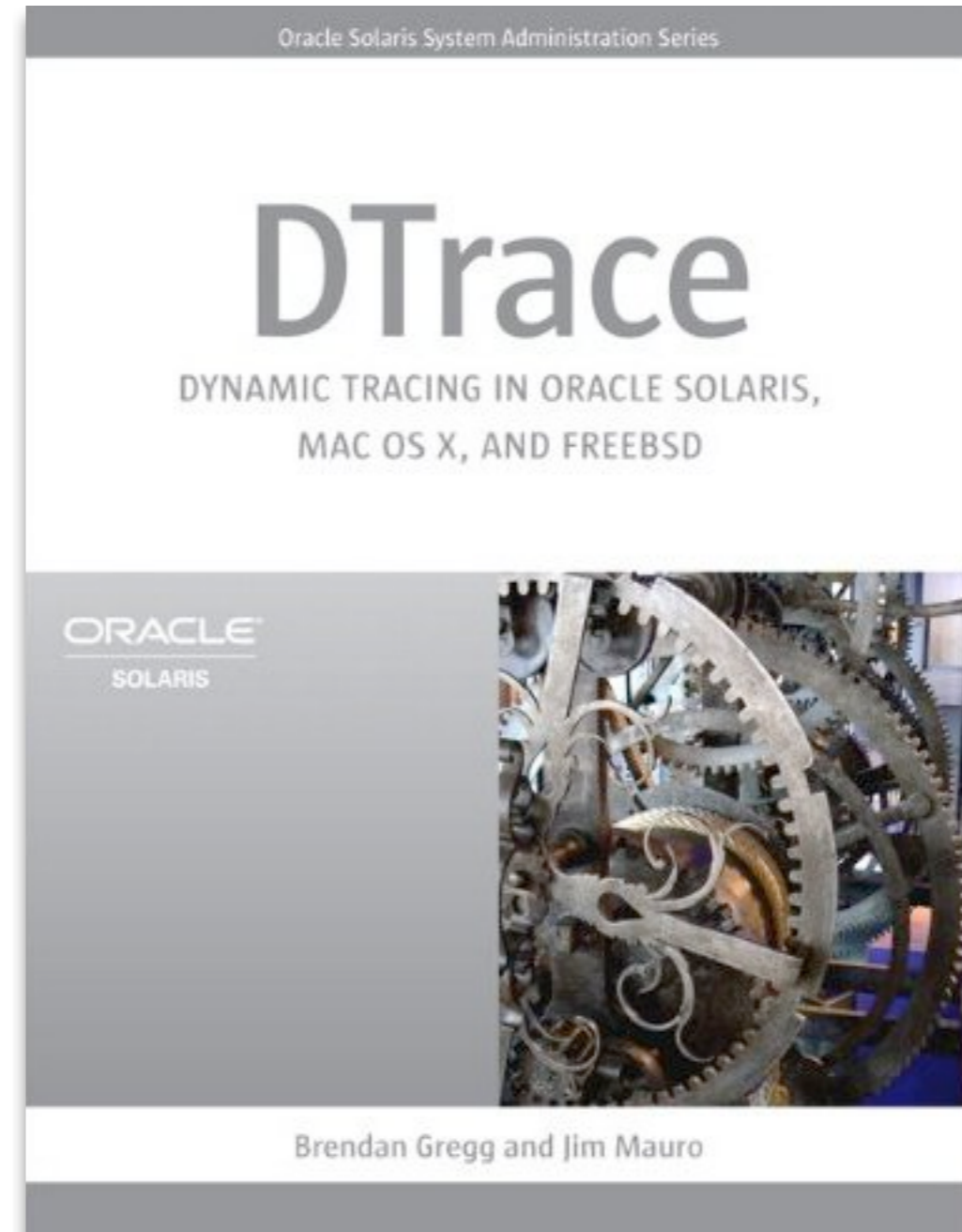
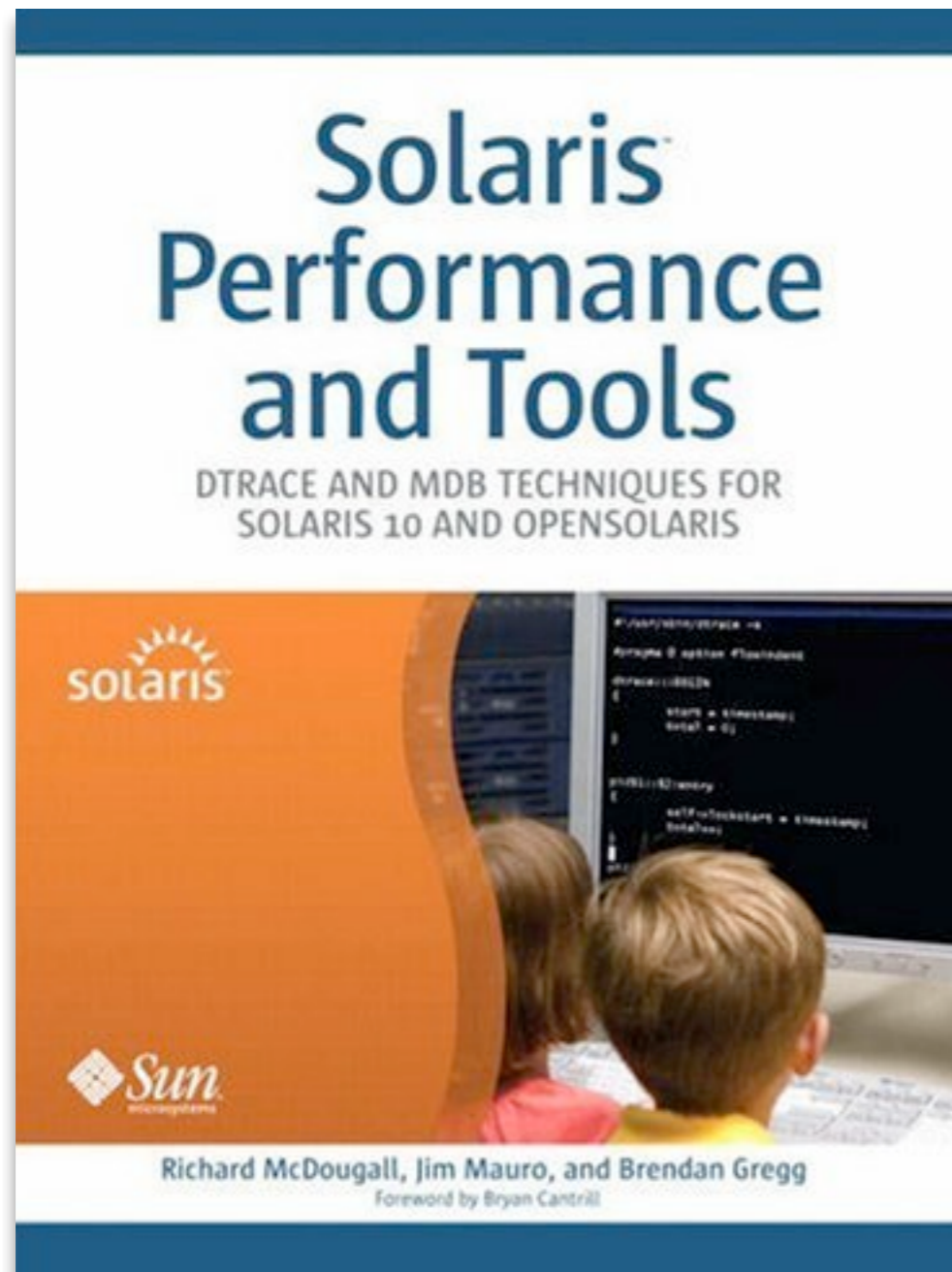
由 Flow Yu 於 1 year 之前發表

文章分類: 儲存裝置



這可不是開玩笑來著，而是經過實驗證明的唷！來自 Sun Fishworks 團隊的 Brendan Gregg，或許是在機房裡待久了，於是突發異想跑去對機房中的磁碟陣列大吼大叫一陣，結果就這麼發現了一個真理，那就是：硬碟就跟員工一樣，吼叫並不會讓他們的效率變高，反而還會讓它們心生不爽而降低士氣，私底下就開始搞罷工。

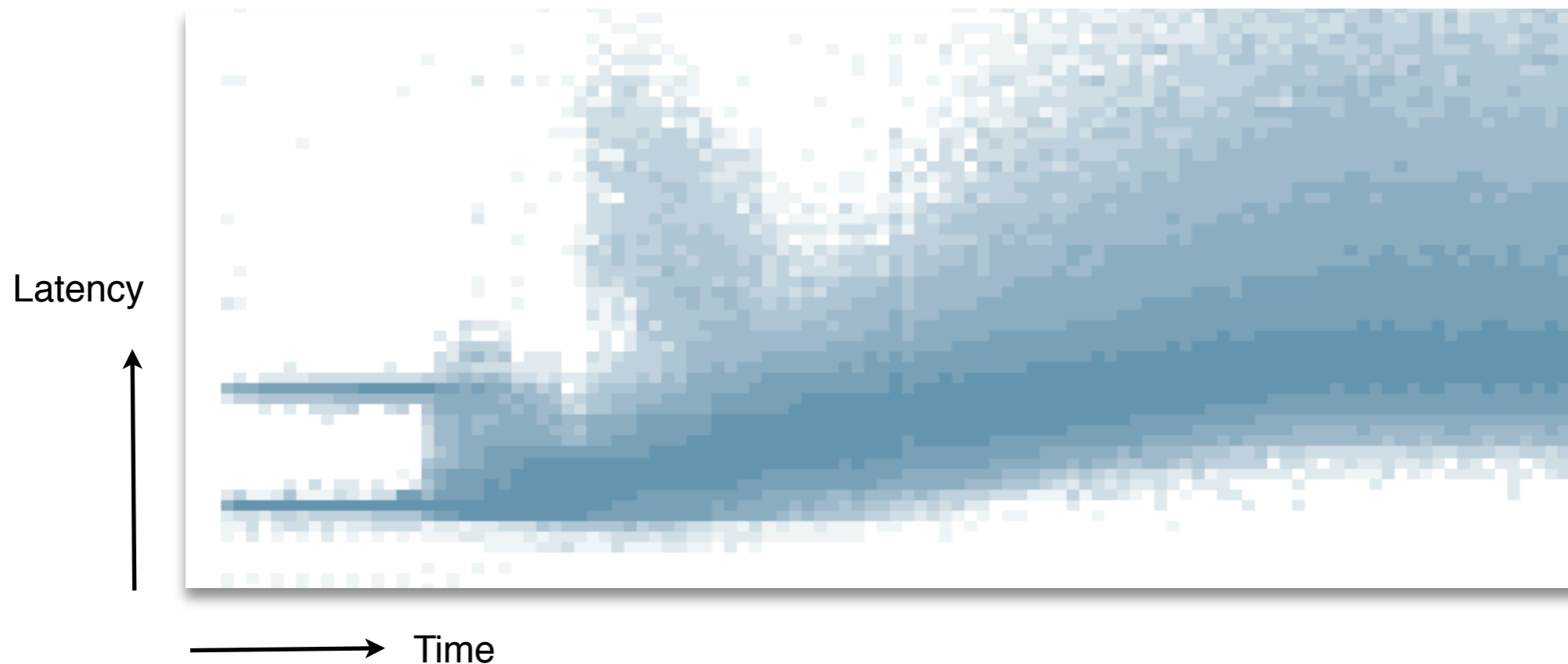
**I do performance analysis
and I'm a DTrace addict**



- Performance
 - Workload Analysis and Resource Monitoring
 - Understanding available and ideal metrics before plotting
- Visualizations
 - Current examples
 - Latency
 - Utilization
 - Future opportunities
 - Cloud Computing

Visualizations like these

- The “rainbow pterodactyl”



- ... which needs quite a bit of explanation

- Consider performance metrics before plotting
- See the value of visualizations
- Remember key examples

- Consider performance metrics before plotting
 - Why studying latency is good
 - ... and studying IOPS can be bad
- See the value of visualizations
 - Why heat maps are needed
 - ... and line graphs can be bad
- Remember key examples
 - I/O latency, as a heat map
 - CPU utilization by CPU, as a heat map

- “Visualizing System Latency”, Communications of the ACM July 2010, by Brendan Gregg
- and more

Understanding the metrics before we visualize them



- Workload analysis
 - Is there an issue? Is an issue real?
 - Where is the issue?
 - Will the proposed fix work? Did it work?
- Resource monitoring
 - How utilized are the environment components?
 - Important activity for capacity planning

- Applied during:
 - software and hardware development
 - proof of concept testing
 - regression testing
 - benchmarking
 - monitoring

- Load

- Architecture

- Load
 - Workload applied
 - Too much for the system?
 - Poorly constructed?
- Architecture
 - System configuration
 - Software and hardware bugs

- Identify or confirm if a workload has a performance issue
 - Quantify
- Locate issue
 - Quantify
- Determine, apply and verify solution
 - Quantify

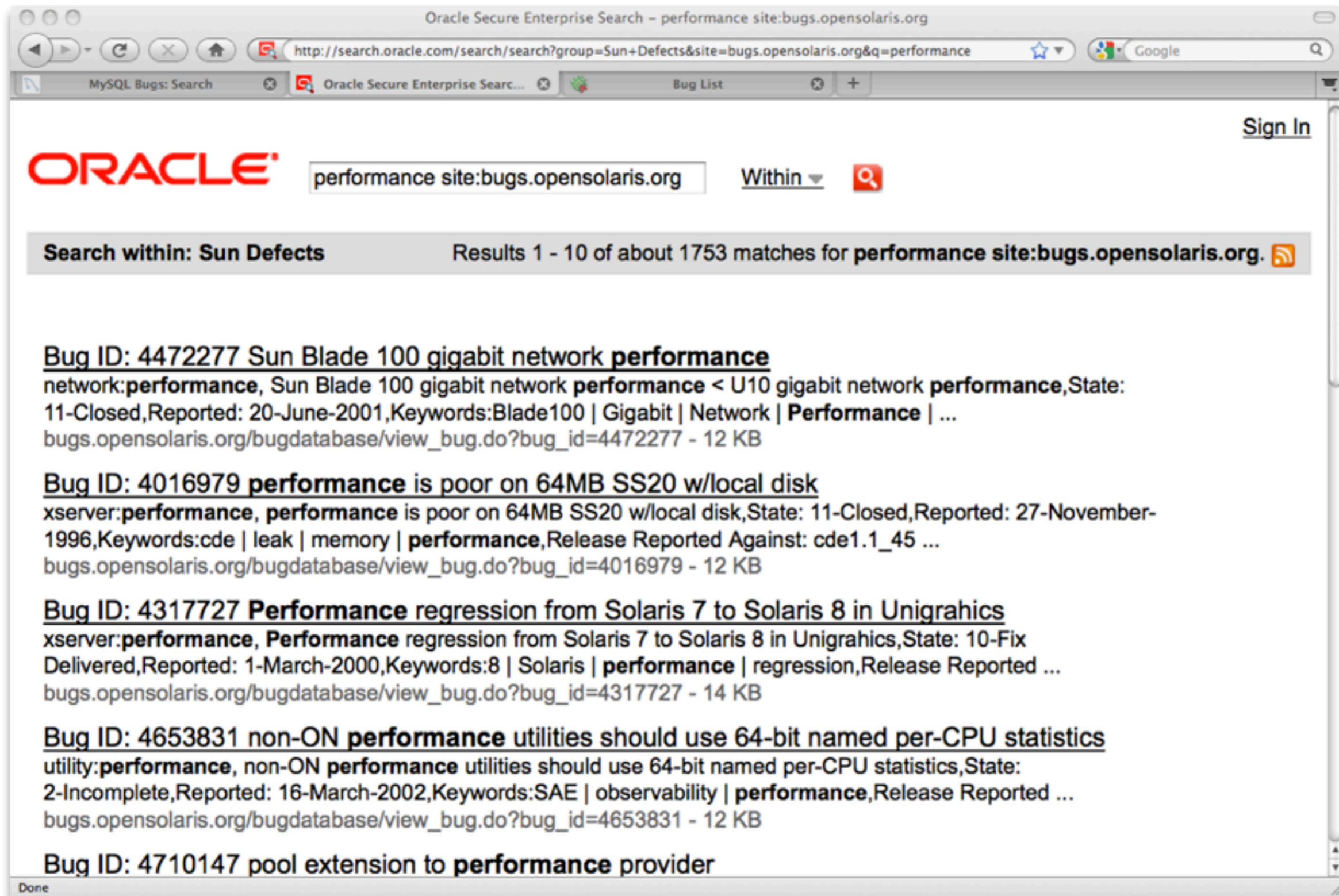
- Finding a performance issue isn't the problem ... it's finding the issue that matters

bugs.mysql.com “performance”



The screenshot shows a web browser window with the URL `http://bugs.mysql.com/search.php?search_for=performance&status=Active&severity=&limit=10&order_by=&cmd`. The page displays a search results table for MySQL bugs related to performance. The table has columns for ID#, Date, Type, Status, Sev, Version, Target, OS, Summary, and Assigned. Five bugs are visible in the list.

ID#	Date	Type	Status	Sev	Version	Target	OS	Summary	Assigned
53696	2010-05-17 13:07	Server: Performance Schema	Verified (36 days)	S3	5.6.99	5.5+	Any	Performance schema engine violates the PSEA API by calling <code>my_error()</code>	Marc Alff
46886	2009-08-24 12:44	Tests: Cluster	Open (397 days)	S3			Any	Make the cluster regression performance testing more stable	Jørgen Austvik
48767	2009-11-13 21:17	Server: DB2SE for IBM i	Open	S5			Any	IBMDB2I subselect performance degrades when async buffering enabled	Tim Clark
37703	2008-06-27 23:20	Server: General	Verified (241 days)	S5	MYSQL 6.0.4 (Source Distribution)		Linux (EL5.1)	MySQL performance with and without Fast Mutexes using Sysbench Workload	
57012	2010-09-25	MySQL Workbench:	Verified	S5	5.2.28	wb53	Linux	Workbench on Linux doesn't use GPU to	



Oracle Secure Enterprise Search – performance site:bugs.opensolaris.org

http://search.oracle.com/search/search?group=Sun+Defects&site=bugs.opensolaris.org&q=performance

MySQL Bugs: Search Oracle Secure Enterprise Search Bug List

Sign In

ORACLE performance site:bugs.opensolaris.org Within

Search within: Sun Defects Results 1 - 10 of about 1753 matches for performance site:bugs.opensolaris.org

Bug ID: 4472277 Sun Blade 100 gigabit network performance
network:performance, Sun Blade 100 gigabit network performance < U10 gigabit network performance,State: 11-Closed,Reported: 20-June-2001,Keywords:Blade100 | Gigabit | Network | Performance | ...
bugs.opensolaris.org/bugdatabase/view_bug.do?bug_id=4472277 - 12 KB

Bug ID: 4016979 performance is poor on 64MB SS20 w/local disk
xserver:performance, performance is poor on 64MB SS20 w/local disk,State: 11-Closed,Reported: 27-November-1996,Keywords:cde | leak | memory | performance,Release Reported Against: cde1.1_45 ...
bugs.opensolaris.org/bugdatabase/view_bug.do?bug_id=4016979 - 12 KB

Bug ID: 4317727 Performance regression from Solaris 7 to Solaris 8 in Unigrahics
xserver:performance, Performance regression from Solaris 7 to Solaris 8 in Unigrahics,State: 10-Fix Delivered,Reported: 1-March-2000,Keywords:8 | Solaris | performance | regression,Release Reported ...
bugs.opensolaris.org/bugdatabase/view_bug.do?bug_id=4317727 - 14 KB

Bug ID: 4653831 non-ON performance utilities should use 64-bit named per-CPU statistics
utility:performance, non-ON performance utilities should use 64-bit named per-CPU statistics,State: 2-Incomplete,Reported: 16-March-2002,Keywords:SAE | observability | performance,Release Reported ...
bugs.opensolaris.org/bugdatabase/view_bug.do?bug_id=4653831 - 12 KB

Bug ID: 4710147 pool extension to performance provider

Done

bugs.mozilla.org: “performance”



Bugzilla@Mozilla - Bug List

Home | New | Browse | Search | performance Search [?] | Reports | Requests | Help | New Account | Log In | Forgot Password

Mon Nov 8 2010 22:45:12 PST

This list is too long for Bugzilla's little mind; the Next/Prev/First/Last buttons won't appear on individual bugs.

Status: REOPENED, NEW, ASSIGNED, UNCONFIRMED **Product:** performance **Component:** performance **Alias:** performance
Summary: performance **Whiteboard:** performance **Content:** "performance"

1995 bugs found.

ID	Sev	Pri	OS	Assignee	Status	Resolution	Summary
491602	nor	--	All	bruce	UNCO	---	Investigate JSGC performance on real workload
466515	nor	--	Wind	create-and-change	UNCO	---	"connection interrupted" received after "Change My Votes"
501515	nor	--	Mac	dmandelin	UNCO	---	Performance bottleneck for object creation in js_NewGCThir
598466	nor	--	Wind	dsicore	UNCO	---	Increased of memory usage between FF3.6 and FF4b6
371063	min	--	Mac	general	UNCO	---	Mouse wheel is not working with drop down menu list
398546	nor	--	All	general	UNCO	---	/config.cgi?ctype=rdf produces XML with about 30% whites bandwidth
412373	nor	--	All	general	UNCO	---	String assignment too slow
458749	enh	--	All	general	UNCO	---	patch to fire interruptHandler once per line of script

“performance” bugs

- ... and those are just the known performance bugs
- ... and usually only of a certain type (architecture)

- Observation based
 - Choose a reliable metric
 - Estimate performance gain from resolving issue
- Experimentation based
 - Apply fix
 - Quantify before vs. after using a reliable metric

- For example:
 - Observed: application I/O takes 10 ms
 - Observed: 9 ms of which is disk I/O
 - Suggestion: replace disks with flash-memory based SSDs, with an expected latency of ~ 100 us
 - Estimated gain: 10 ms \rightarrow 1.1 ms (10 ms - 9 ms + 0.1 ms)
 $\approx \sim 9$ x gain
- Very useful - but not possible without accurate quantification

- For example:
 - Observed: Application transaction latency average 10 ms
 - Experiment: Added more DRAM to increase cache hits and reduce average latency
 - Observed: Application transaction latency average 2 ms
 - Gain: 10 ms \rightarrow 2 ms = 5x

- Also very useful - but risky without accurate quantification

- Choose reliable metrics to quantify performance:
 - IOPS
 - transactions/second
 - throughput
 - utilization
 - latency
- Ideally
 - interpretation is straightforward
 - reliable

- Choose reliable metrics to quantify performance:

- IOPS
- transactions/second
- throughput
- utilization
- latency

generally better suited for:

Capacity Planning

Workload Analysis

- Ideally

- interpretation is straightforward
- reliable

- Ideally (given the luxury of time):
 - design the desired metrics
 - then see if they exist, or,
 - implement them (eg, DTrace)
- Non-ideally
 - see what already exists
 - make-do (eg, vmstat -> gnuplot)

- Available metrics are implemented correctly
 - all software has bugs
 - eg, CR: 6687884 `nxge rbytes` and `obytes kstat` are wrong
 - trust no metric without double checking from other sources
- Available metrics are designed by performance experts
 - sometimes added by the programmer to only debug their code
- Available metrics are complete
 - you won't always find what you really need

- This will be explained using two examples:
 - Workload Analysis
 - Capacity Planning

- Quantifying performance issues with IOPS vs latency
 - IOPS is commonly presented by performance analysis tools
 - eg: disk IOPS via kstat, SNMP, iostat, ...

- Depends on where the I/O is measured
 - app -> library -> syscall -> VFS -> filesystem -> RAID -> device
- Depends on what the I/O is
 - synchronous or asynchronous
 - random or sequential
 - size
- Interpretation difficult
 - what value is good or bad?
 - is there a max?

- IOPS Inflation
 - Library or Filesystem prefetch/read-ahead
 - Filesystem metadata
 - RAID stripes
- IOPS Deflation
 - Read caching
 - Write cancellation
 - Filesystem I/O aggregation
- IOPS aren't created equal

IOPS example: iostat -xnz 1



- Consider this disk: 86 IOPS == 99% busy

```
extended device statistics
  r/s    w/s    kr/s    kw/s  wait  actv  wsvc_t  asvc_t   %w   %b  device
 86.6    0.0   655.5    0.0   0.0   1.0    0.0    11.5     0   99  c1d0
```

- Versus this disk: 21,284 IOPS == 99% busy

```
extended device statistics
  r/s    w/s    kr/s    kw/s  wait  actv  wsvc_t  asvc_t   %w   %b  device
21284.7  0.0 10642.4    0.0   0.0   1.8    0.0    0.1     2   99  c1d0
```

IOPS example: iostat -xnz 1



- Consider this disk: 86 IOPS == 99% busy

```
extended device statistics
  r/s    w/s    kr/s    kw/s  wait  actv  wsvc_t  asvc_t  %w  %b  device
 86.6    0.0   655.5    0.0   0.0   1.0    0.0    11.5    0   99  c1d0
```

- Versus this disk: 21,284 IOPS == 99% busy

```
extended device statistics
  r/s    w/s    kr/s    kw/s  wait  actv  wsvc_t  asvc_t  %w  %b  device
21284.7  0.0  10642.4  0.0   0.0   1.8    0.0    0.1    2   99  c1d0
```

- ... they are the same disk, different I/O types
 - 1) 8 Kbyte random
 - 2) 512 byte sequential (on-disk DRAM cache)

Using IOPS to quantify issues



- to identify
 - is 100 IOPS an problem? Per disk?
- to locate
 - 90% of IOPS are random. Is that the problem?
- to verify
 - A filesystem tunable caused IOPS to reduce. Has this fixed the issue?

- to identify
 - is 100 IOPS an problem? Per disk? (depends...)
- to locate
 - 90% of IOPS are random. Is that the problem? (depends...)
- to verify
 - A filesystem tunable caused IOPS to reduce.
Has this fixed the issue? (probably, assuming...)
- We can introduce more metrics to understand these, but standalone IOPS is tricky to interpret

Using latency to quantify issues



- to identify
 - is a 100ms I/O a problem?
- to locate
 - 90ms of the 100ms is lock contention. Is that the problem?
- to verify
 - A filesystem tunable caused the I/O latency to reduce to 1ms. Has this fixed the issue?

- to identify
 - is a 100ms I/O a problem? (probably - if synchronous to the app.)
- to locate
 - 90ms of the 100ms is lock contention. Is that the problem? (yes)
- to verify
 - A filesystem tunable caused the I/O latency to reduce to 1ms. Has this fixed the issue? (probably - if 1ms is acceptable)
- Latency is much more reliable, easier to interpret

- Time from I/O or transaction request to completion
- Synchronous latency has a direct impact on performance
 - Application is waiting
 - higher latency == worse performance
- Not all latency is synchronous:
 - Asynchronous filesystem threads flushing dirty buffers to disk
eg, zfs TXG synchronous thread
 - Filesystem prefetch
no one is waiting at this point
 - TCP buffer and congestion window: individual packet latency may be high, but pipe is kept full for good throughput performance

- Currency converter (* -> ms):
 - random disk IOPS == I/O service latency
 - disk saturation == I/O wait queue latency
 - CPU utilization == code path execution latency
 - CPU saturation == dispatcher queue latency
 - ...
- Quantifying as latency allows different components to be compared, ratios examined, improvements estimated.

- Different performance activity
 - Focus is environment components, not specific issues
 - incl. CPUs, disks, network interfaces, memory, I/O bus, memory bus, CPU interconnect, I/O cards, network switches, etc.
 - Information is used for capacity planning
 - Identifying future issues before they happen
- Quantifying resource monitoring with IOPS vs utilization

- Another look at this disk:

```
extended device statistics
  r/s    w/s    kr/s    kw/s  wait  actv  wsvc_t  asvc_t  %w  %b  device
  86.6    0.0   655.5    0.0   0.0   1.0    0.0    11.5    0   99  c1d0
[...]
```

```
extended device statistics
  r/s    w/s    kr/s    kw/s  wait  actv  wsvc_t  asvc_t  %w  %b  device
 21284.7  0.0 10642.4    0.0   0.0   1.8    0.0    0.1    2   99  c1d0
```

- Q. does this system need more spindles for IOPS capacity?

- Another look at this disk:

```
extended device statistics
  r/s    w/s    kr/s    kw/s  wait  actv  wsvc_t  asvc_t  %w  %b  device
  86.6   0.0   655.5   0.0   0.0   1.0   0.0     11.5    0  99  c1d0
[...]
```

```
extended device statistics
  r/s    w/s    kr/s    kw/s  wait  actv  wsvc_t  asvc_t  %w  %b  device
21284.7  0.0 10642.4  0.0   0.0   1.8   0.0     0.1    2  99  c1d0
```

- Q. does this system need more spindles for IOPS capacity?
 - IOPS (r/s + w/s): ???
 - Utilization (%b): yes (even considering NCQ)

- Another look at this disk:

```
extended device statistics
  r/s    w/s    kr/s    kw/s  wait  actv  wsvc_t  asvc_t  %w  %b  device
  86.6   0.0   655.5   0.0   0.0   1.0   0.0     11.5    0  99  c1d0
[...]
```

```
extended device statistics
  r/s    w/s    kr/s    kw/s  wait  actv  wsvc_t  asvc_t  %w  %b  device
 21284.7  0.0 10642.4  0.0   0.0   1.8   0.0     0.1    2  99  c1d0
```

- Q. does this system need more spindles for IOPS capacity?
 - IOPS (r/s + w/s): ???
 - Utilization (%b): yes (even considering NCQ)
 - Latency (wsvc_t): no
- Latency will identify the issue once it is an issue; utilization will forecast the issue - capacity planning

- Metrics matter - need to reliably quantify performance
 - to identify, locate, verify
 - try to think, design
- Workload Analysis
 - latency
- Resource Monitoring
 - utilization
- Other metrics are useful to further understand the nature of the workload and resource behavior

- Consider performance metrics before plotting
 - Why latency is good
 - ... and IOPS can be bad
- See the value of visualizations
 - Why heat maps are needed
 - ... and line graphs can be bad
- Remember key examples
 - I/O latency, as a heat map
 - CPU utilization by CPU, as a heat map



Current Examples

Latency



- So far we've picked:
- Latency
 - for workload analysis
- Utilization
 - for resource monitoring

- For example, disk I/O
- Raw data looks like this:

```
# iosnoop -o
DTIME      UID    PID D    BLOCK  SIZE    COMM  PATHNAME
125        100    337 R    72608  8192    bash  /usr/sbin/tar
138        100    337 R    72624  8192    bash  /usr/sbin/tar
127        100    337 R    72640  8192    bash  /usr/sbin/tar
135        100    337 R    72656  8192    bash  /usr/sbin/tar
118        100    337 R    72672  8192    bash  /usr/sbin/tar
108        100    337 R    72688  4096    bash  /usr/sbin/tar
87         100    337 R    72696  3072    bash  /usr/sbin/tar
9148       100    337 R    113408 8192    tar   /etc/default/lu
8806       100    337 R    104738 7168    tar   /etc/default/lu
2262       100    337 R    13600  1024    tar   /etc/default/cron
76         100    337 R    13616  1024    tar   /etc/default/devfsadm
```

[...many pages of output...]

- iosnoop is DTrace based
 - examines latency for every disk (back end) I/O

- tuples
 - I/O completion time
 - I/O latency
- can be 1,000s of these per second

Summarizing Latency



- iostat(1M) can show per second average:

```
$ iostat -xnz 1
[...]
```

extended device statistics										
r/s	w/s	kr/s	kw/s	wait	actv	wsvc_t	asvc_t	%w	%b	device
471.0	7.0	786.1	12.0	0.1	1.2	0.2	2.5	4	90	c1d0
extended device statistics										
r/s	w/s	kr/s	kw/s	wait	actv	wsvc_t	asvc_t	%w	%b	device
631.0	0.0	1063.1	0.0	0.2	1.0	0.3	1.6	9	92	c1d0
extended device statistics										
r/s	w/s	kr/s	kw/s	wait	actv	wsvc_t	asvc_t	%w	%b	device
472.0	0.0	529.0	0.0	0.0	1.0	0.0	2.1	0	94	c1d0

```
[...]
```

- Condenses I/O completion time
- Almost always a sufficient resolution
 - (So far I've only had one case where examining raw completion time data was crucial: an interrupt coalescing bug)

- Average loses latency outliers
- Average loses latency distribution
- ... but not disk distribution:

```
$ iostat -xnz 1  
[...]
```

```
                extended device statistics  
   r/s    w/s    kr/s    kw/s  wait  actv  wsvc_t  asvc_t  %w  %b  device  
   43.9    0.0   351.5    0.0   0.0   0.4    0.0    10.0    0  34  c0t5000CCA215C46459d0  
   47.6    0.0   381.1    0.0   0.0   0.5    0.0     9.8    0  36  c0t5000CCA215C4521Dd0  
   42.7    0.0   349.9    0.0   0.0   0.4    0.0    10.1    0  35  c0t5000CCA215C45F89d0  
   41.4    0.0   331.5    0.0   0.0   0.4    0.0     9.6    0  32  c0t5000CCA215C42A4Cd0  
   45.6    0.0   365.1    0.0   0.0   0.4    0.0     9.2    0  34  c0t5000CCA215C45541d0  
   45.0    0.0   360.3    0.0   0.0   0.4    0.0     9.4    0  34  c0t5000CCA215C458F1d0  
   42.9    0.0   343.5    0.0   0.0   0.4    0.0     9.9    0  33  c0t5000CCA215C450E3d0  
   44.9    0.0   359.5    0.0   0.0   0.4    0.0     9.3    0  35  c0t5000CCA215C45323d0  
[...]
```

- only because iostat(1M) prints this per-disk
 - but that gets hard to read for 100s of disks, per second!

- Occasional high-latency I/O
- Can be the sole reason for performance issues
- Can be lost in an average
 - 10,000 fast I/O @ 1ms
 - 1 slow I/O @ 500ms
 - average = 1.05 ms
- Can be seen using max instead of (or as well as) average

- iostat(1M) doesn't show this, however DTrace can
- can be visualized along with average/second
- does identify outliers
- doesn't identify latency distribution details

- Apart from outliers and average, it can be useful to examine the full profile of latency - all the data.
 - For such a crucial metric, keep as much details as possible
- For latency, distributions we'd expect to see include:
 - bi-modal: cache hit vs cache miss
 - tri-modal: multiple cache layers
 - flat: random disk I/O

Latency Distribution Example



- Using DTrace:

```
# ./disklatency.d
Tracing... Hit Ctrl-C to end.
^C
sd4 (28,256), us:
```

value	----- Distribution -----	count
16		0
32		82
64	@@@	621
128	@@@@@	833
256	@@@@@	641
512	@@@	615
1024	@@@@@@@@	1239
2048	@@@@@@@@@@	1615
4096	@@@@@@@@@@	1483
8192		76
16384		1
32768		0
65536		2
131072		0

- not why we are here, but before someone asks...

```
#!/usr/sbin/dtrace -s

#pragma D option quiet

dtrace:::BEGIN
{
    printf("Tracing... Hit Ctrl-C to end.\n");
}

io:::start
{
    start_time[arg0] = timestamp;
}

io:::done
/this->start = start_time[arg0]/
{
    this->delta = (timestamp - this->start) / 1000;
    @[args[1]->dev_statname, args[1]->dev_major, args[1]->dev_minor] =
        quantize(this->delta);
    start_time[arg0] = 0;
}

dtrace:::END
{
    printa("    %s (%d,%d), us:\n%@d\n", @);
}
```

Latency Distribution Example



```
# ./disklatency.d
Tracing... Hit Ctrl-C to end.
^C
sd4 (28,256), us:
```

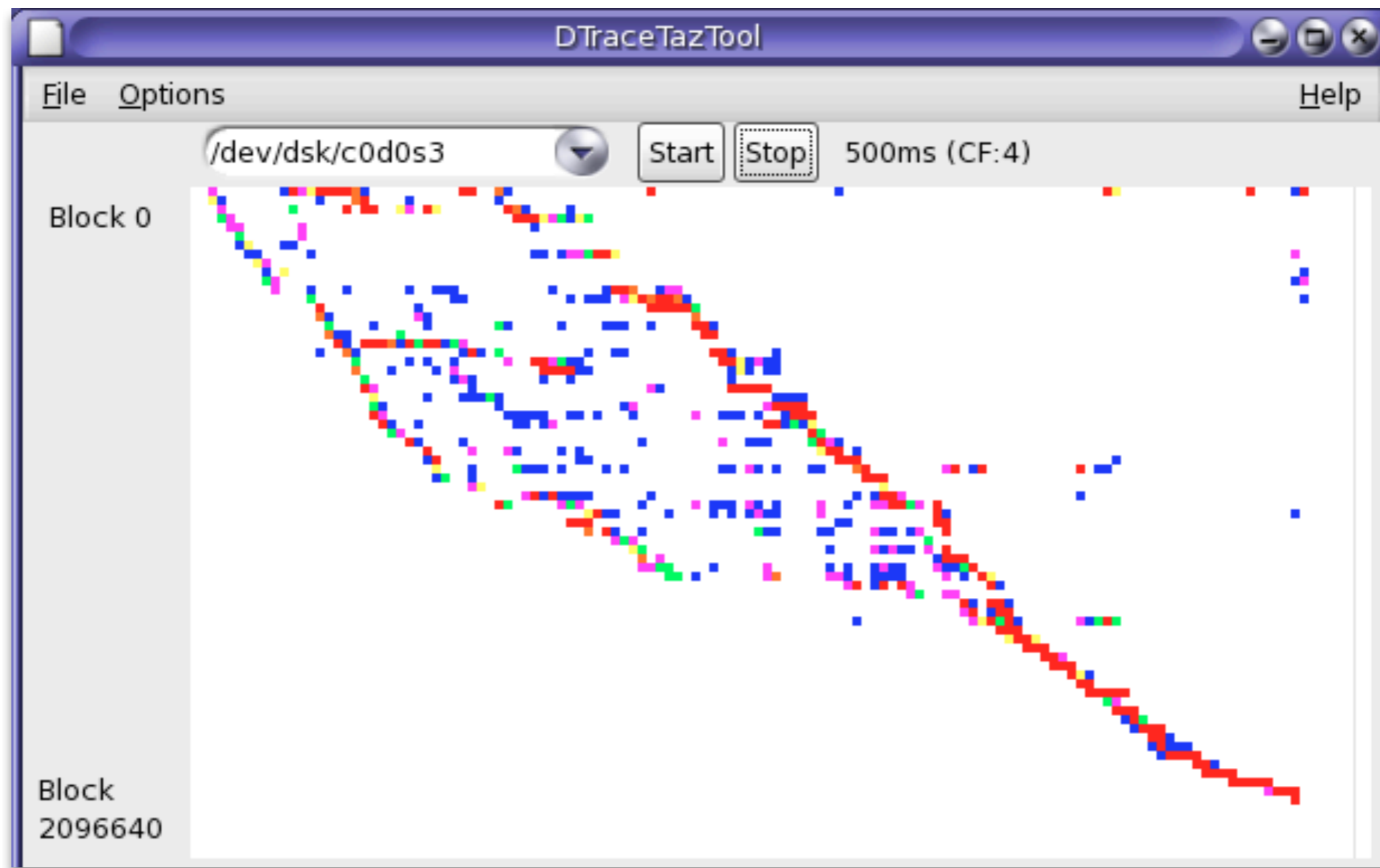
value	----- Distribution -----	count
16		0
32		82
64	@@@	621
128	@@@@@	833
256	@@@@@	641
512	@@@	615
1024	@@@@@@@@	1239
2048	@@@@@@@@@@	1615
4096	@@@@@@@@@@	1483
8192		76
16384		1
32768		0
65536		2
131072		0

→ 65 - 131 ms outliers

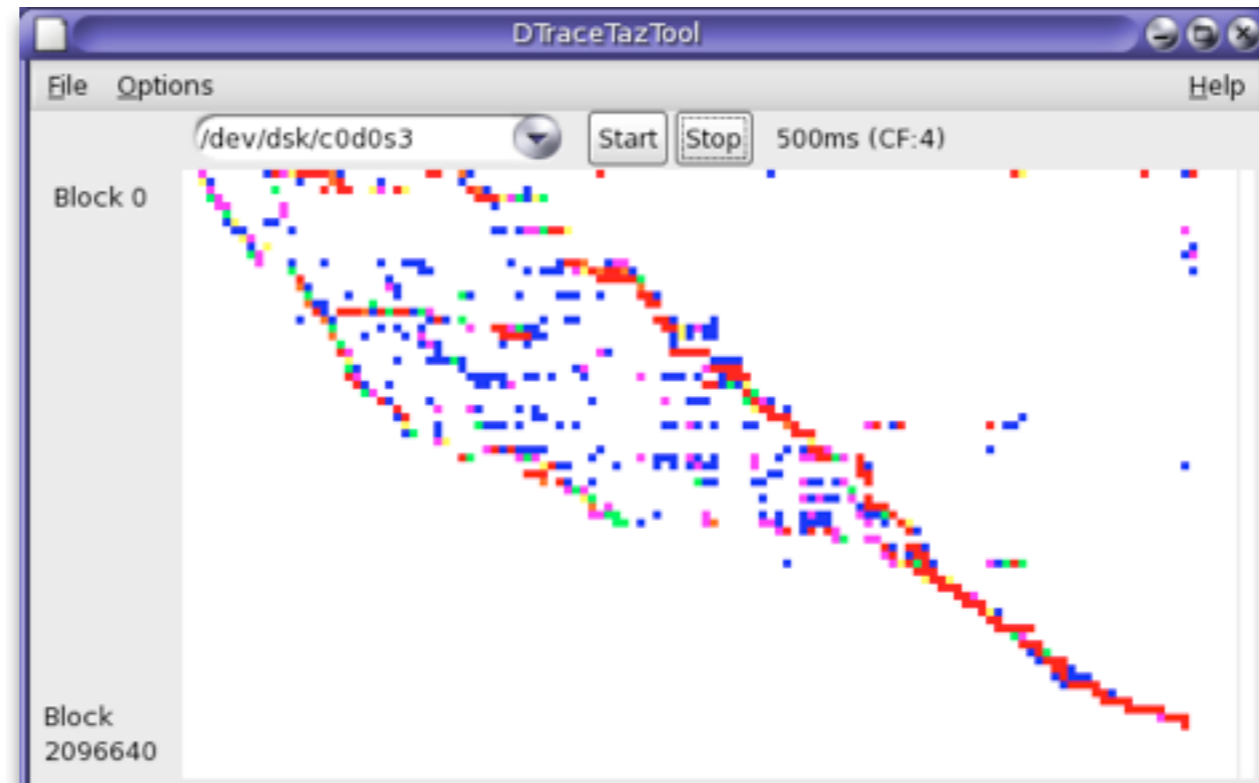
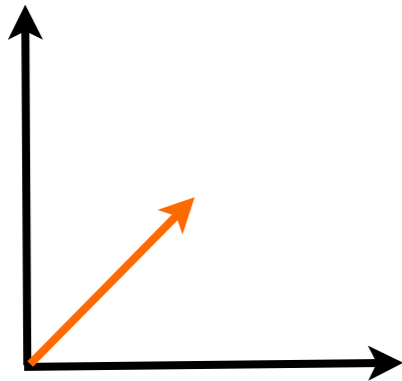
- ... but can we see this distribution per second?
- ... how do we visualize a 3rd dimension?

Column Quantized Visualization aka “heat map”

- For example:



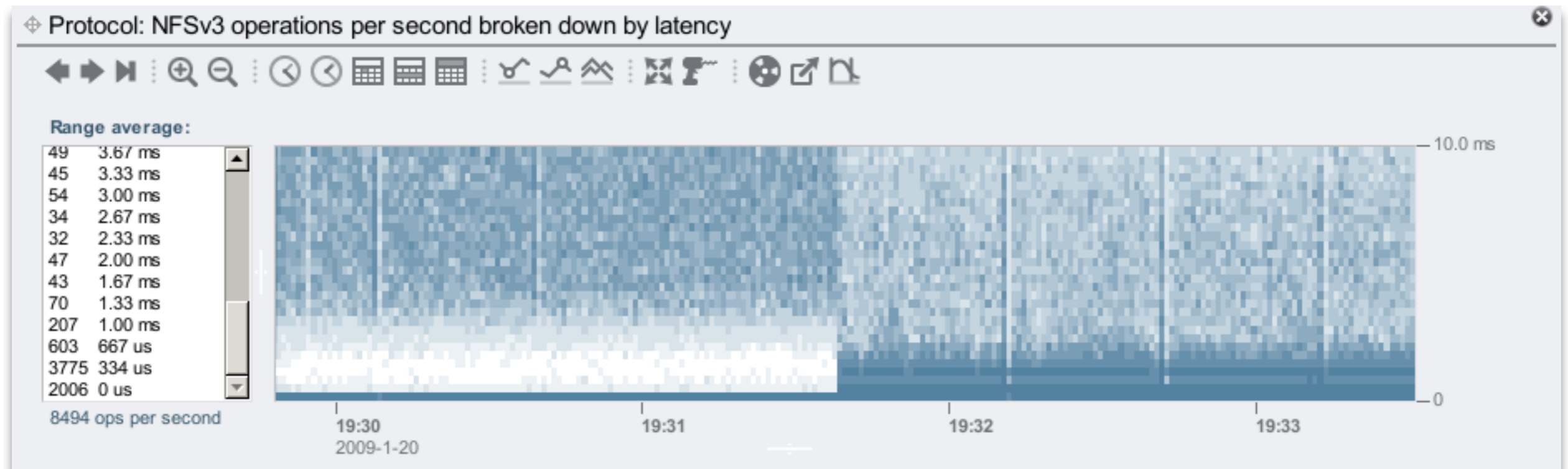
Heat Map: Offset Distribution



- x-axis: time
- y-axis: offset
- z-axis (color scale): I/O count for that time/offset range
- Identified random vs. sequential very well
- Similar heat maps have been used before by defrag tools

Heat Map: Latency Distribution

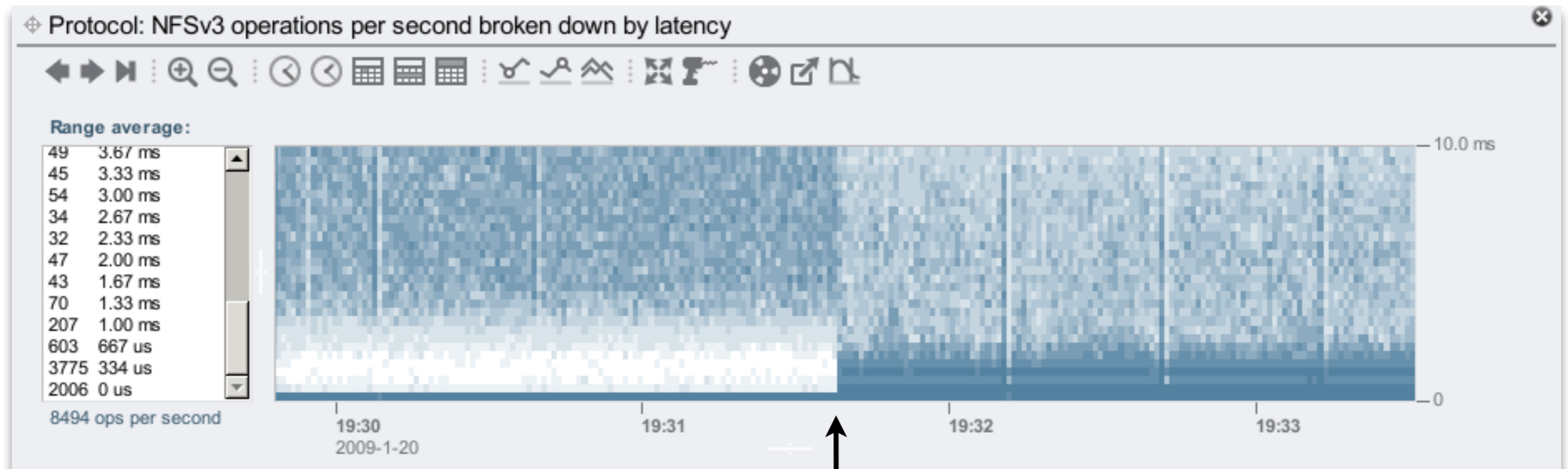
- For example:



- x-axis: time
- y-axis: latency
- z-axis (color saturation): I/O count for that time/latency range

Heat Map: Latency Distribution

- ... in fact, this is a great example:



- reads served from:

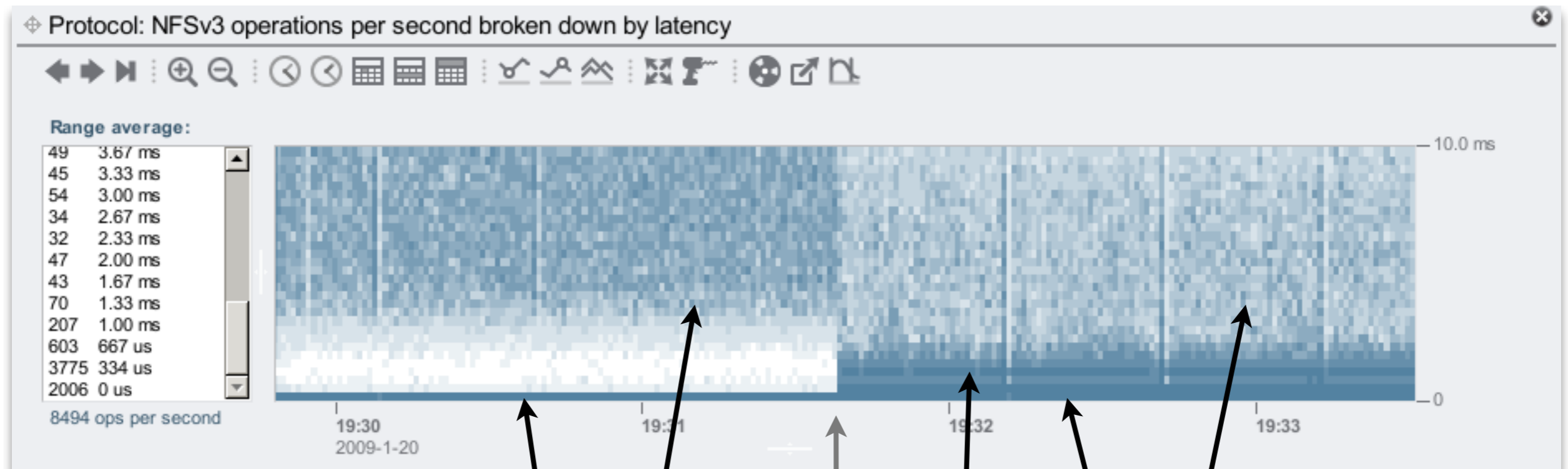
DRAM
disk

DRAM
flash-memory based SSD
disk

ZFS "L2ARC" enabled

Heat Map: Latency Distribution

- ... in fact, this is a great example:



- reads served from:

DRAM
disk

DRAM
flash-memory based SSD
disk

ZFS "L2ARC" enabled

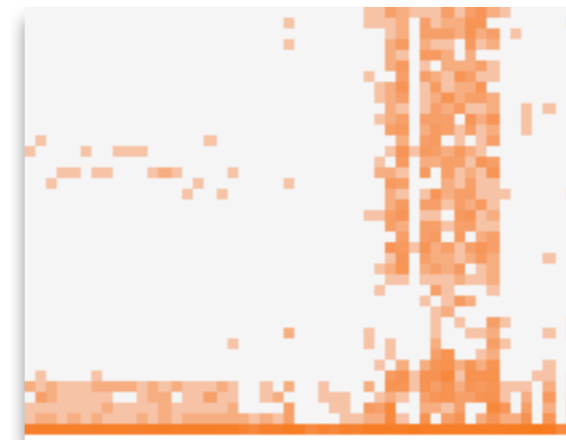
- A color shaded matrix of pixels
- Each pixel is a time and latency range
- Color shade picked based on number of I/O in that range
- Adjusting saturation seems to work better than color hue.

Eg:

- darker == more I/O
- lighter == less I/O

- Large pixels (and corresponding time/latency ranges)

- increases likelihood that adjacent pixels include I/O, have color, and combine to form patterns
- allows color to be more easily seen



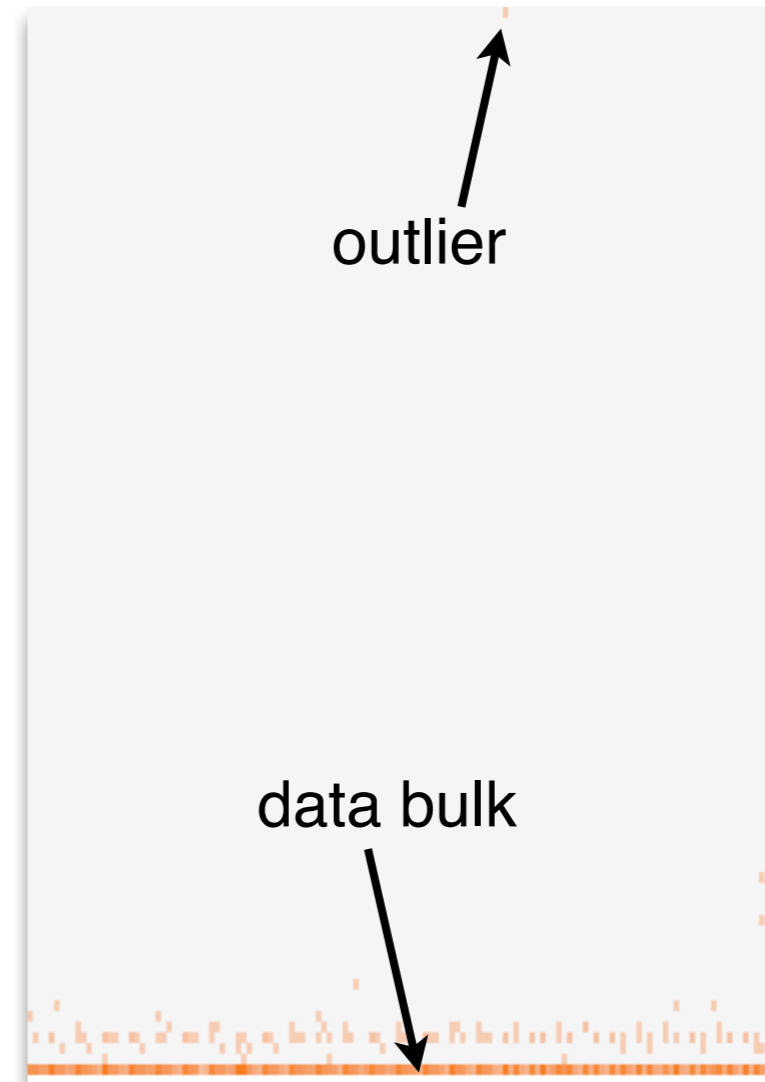
- Smaller pixels (and time/latency ranges)

- can make heat map look like a scatter plot
- of the same color - if ranges are so small only one I/O is typically included



- Linear scale can make subtle details (outliers) difficult to see
 - observing latency outliers is usually of high importance
 - outliers are usually $< 1\%$ of the I/O
 - assigning $< 1\%$ of the color scale to them will washout patterns
- False color palette can be used to emphasize these details
 - although color comparisons become more confusing - non-linear

- Heat maps show these very well
- However, latency outliers can compress the bulk of the heat map data
 - eg, 1 second outlier while most I/O is < 10 ms
- Users should have some control to be able to zoom/truncate details
 - both x and y axis



- Since heat-maps are three dimensions, storing this data can become costly (volume)
- Most of the data points are zero
 - and you can prevent storing zero's by only storing populated elements: associative array
- You can reduce to a sufficiently high resolution, and resample lower as needed
- You can also be aggressive at reducing resolution at higher latencies
 - 10 us granularity not as interesting for I/O > 1 second
 - non-linear resolution

Other Interesting Latency Heat Maps

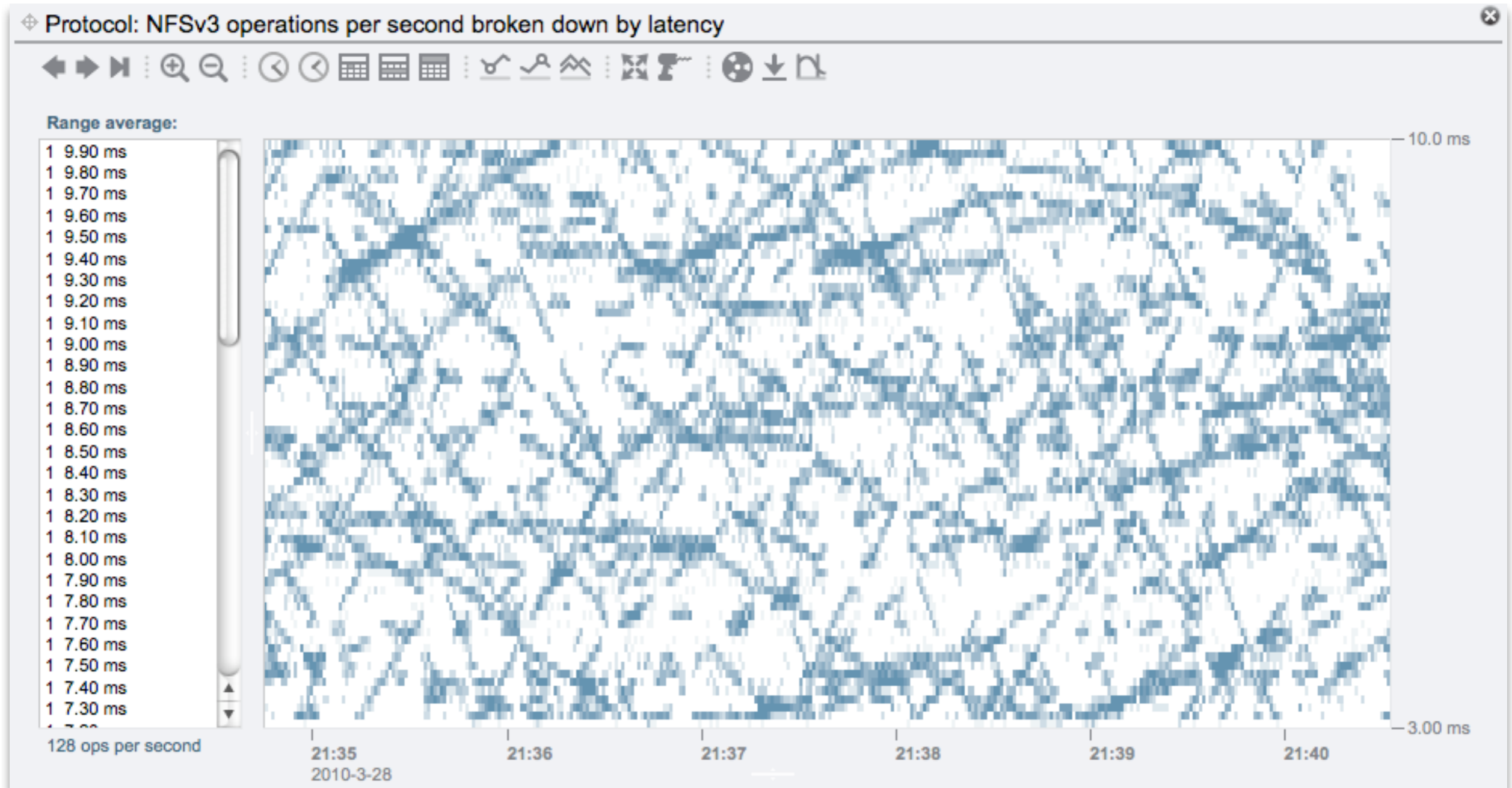


- The “Icy Lake”
- The “Rainbow Pterodactyl”
- Latency Levels

The “Icy Lake” Workload

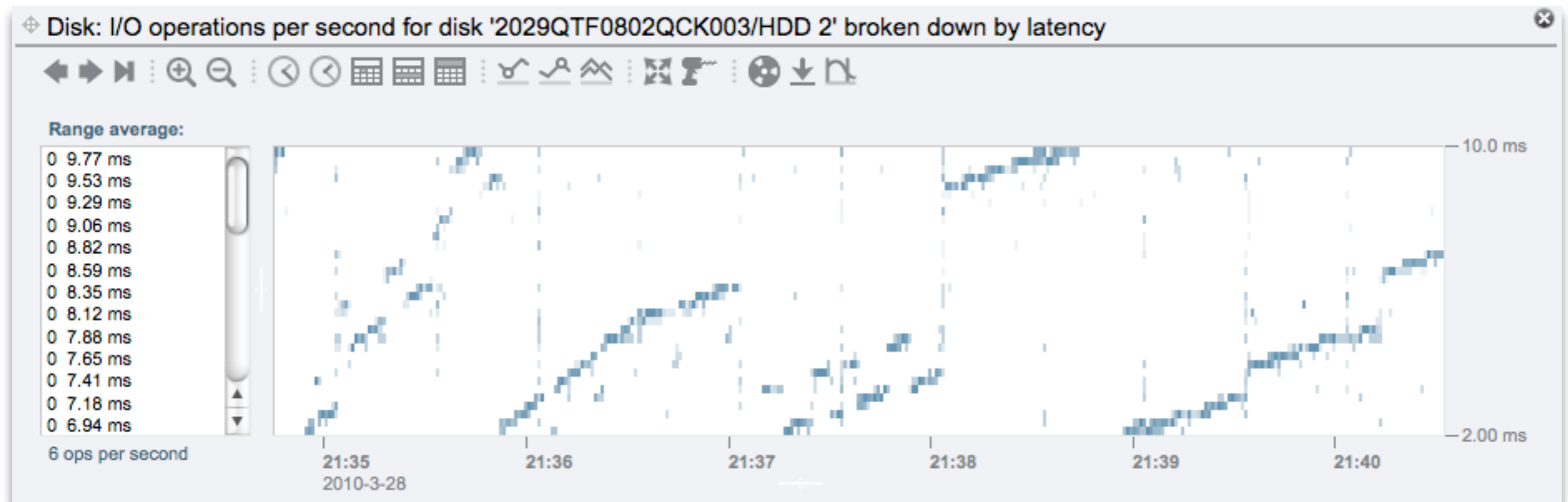
- About as simple as it gets:
 - Single client, single thread, sequential synchronous 8 Kbyte writes to an NFS share
 - NFS server: 22 x 7,200 RPM disks, striped pool
- The resulting latency heat map was unexpected

The “Icy Lake”



“Icy Lake” Analysis: Observation

- Examining single disk latency:

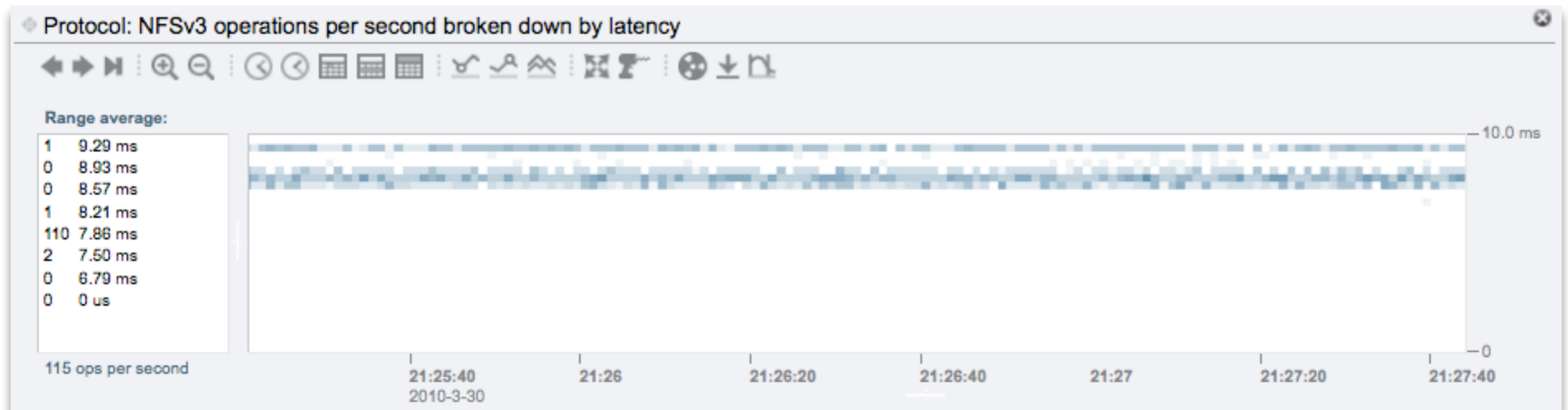


- Pattern match with NFS latency: similar lines
 - each disk contributing some lines to the overall pattern

- We just associated NFS latency with disk device latency, using our eyeballs
 - see the titles on the previous heat maps
- You can programmatically do this (DTrace), but that can get difficult to associate context across software stack layers (but not impossible!)
- Heat Maps allow this part of the problem to be offloaded to your brain
 - and we are very good at pattern matching

“Icy Lake” Analysis: Experimentation

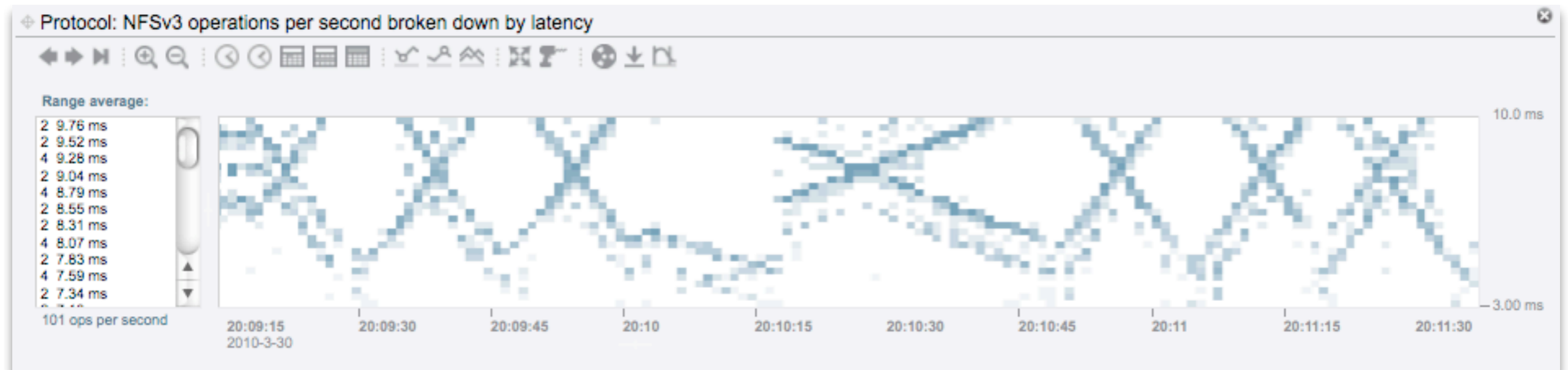
- Same workload, single disk pool:



- No diagonal lines
 - but more questions - see the line (false color palette enhanced) at 9.29 ms? this is < 1% of the I/O. (I'm told, and I believe, that this is due to adjacent track seek latency.)

“Icy Lake” Analysis: Experimentation

- Same workload, two disk striped pool:



- Ah-hah! Diagonal lines.
 - ... but still more questions: why does the angle sometimes change? why do some lines slope upwards and some down?

“Icy Lake” Analysis: Experimentation



- ... each disk from that pool:



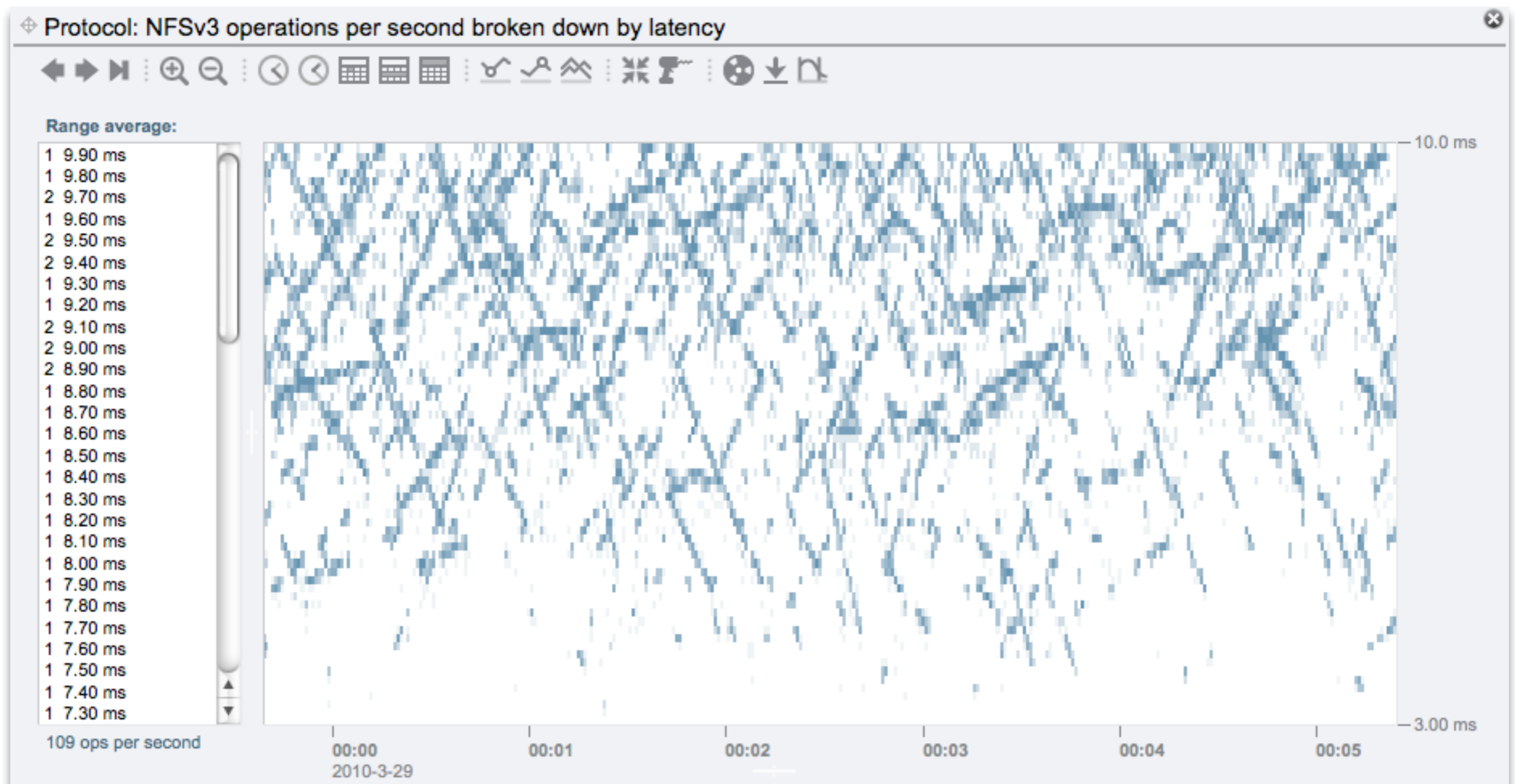
“Icy Lake” Analysis: Questions



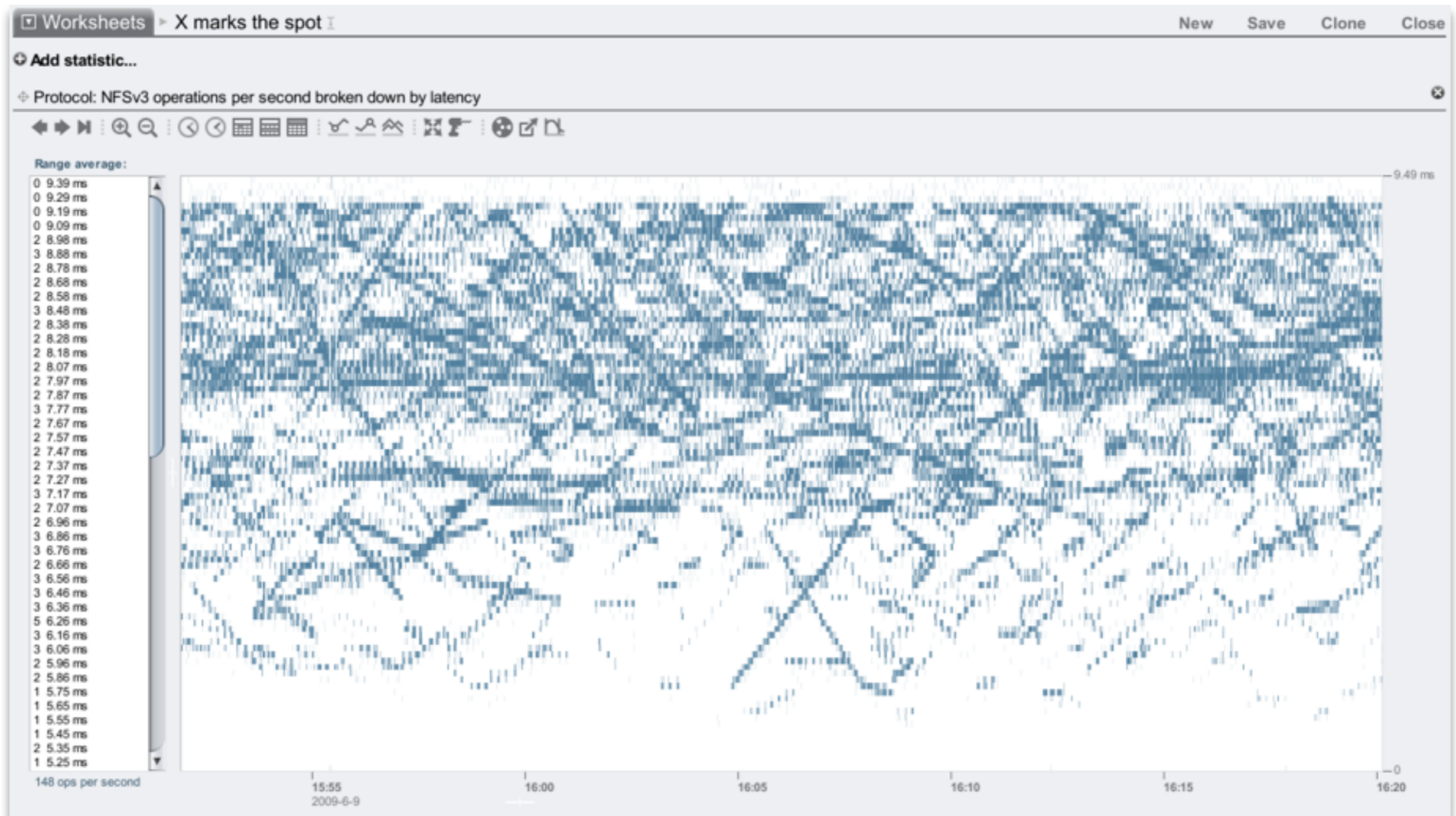
- Remaining Questions:
 - Why does the slope sometimes change?
 - What exactly seeds the slope in the first place?

“Icy Lake” Analysis: Mirroring

- Trying mirroring the pool disks instead of striping:



Another Example: “X marks the spot”

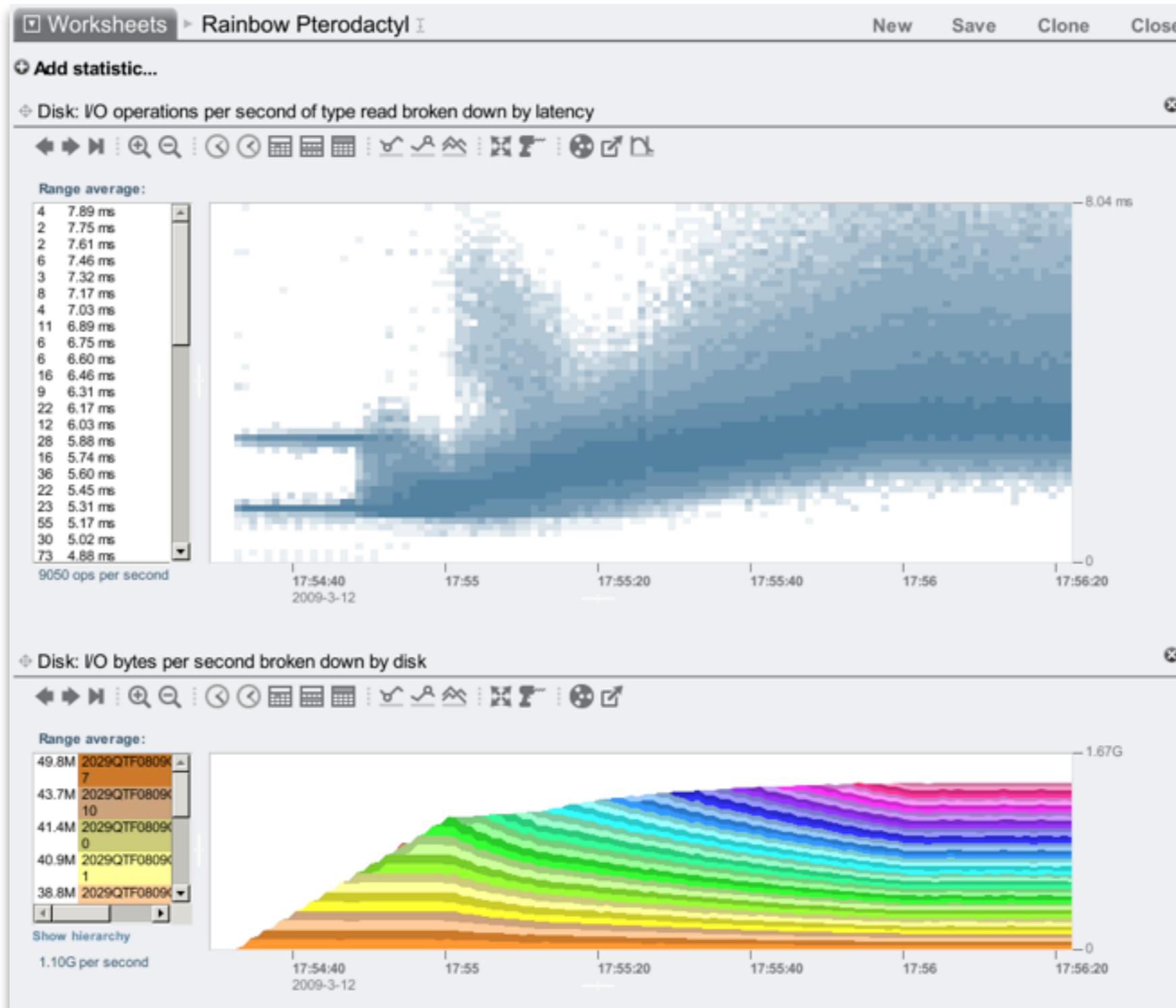


The “Rainbow Pterodactyl” Workload

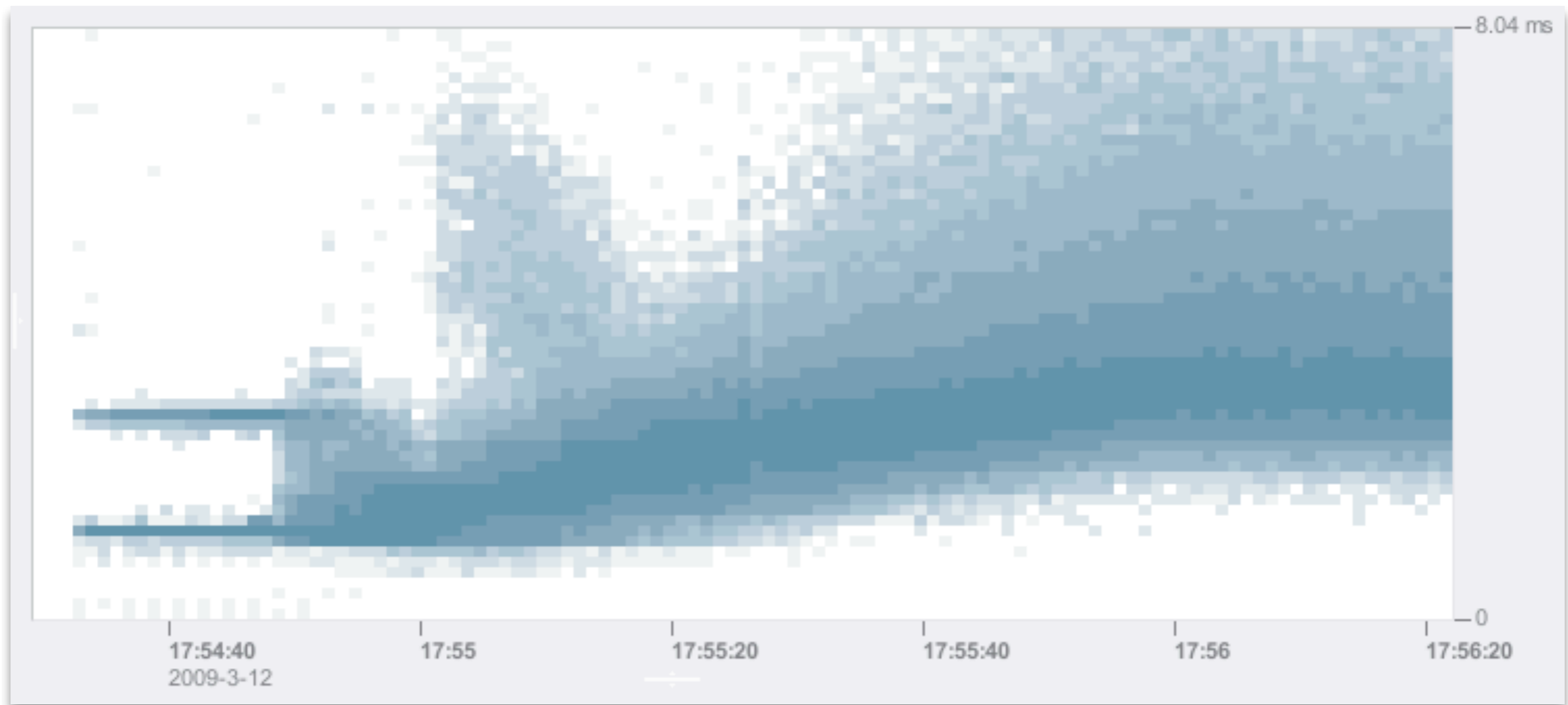


- 48 x 7,200 RPM disks, 2 disk enclosures
- Sequential 128 Kbyte reads to each disk (raw device), adding disks every 2 seconds
- Goal: Performance analysis of system architecture
 - identifying I/O throughput limits by driving I/O subsystem to saturation, one disk at a time (finds knee points)

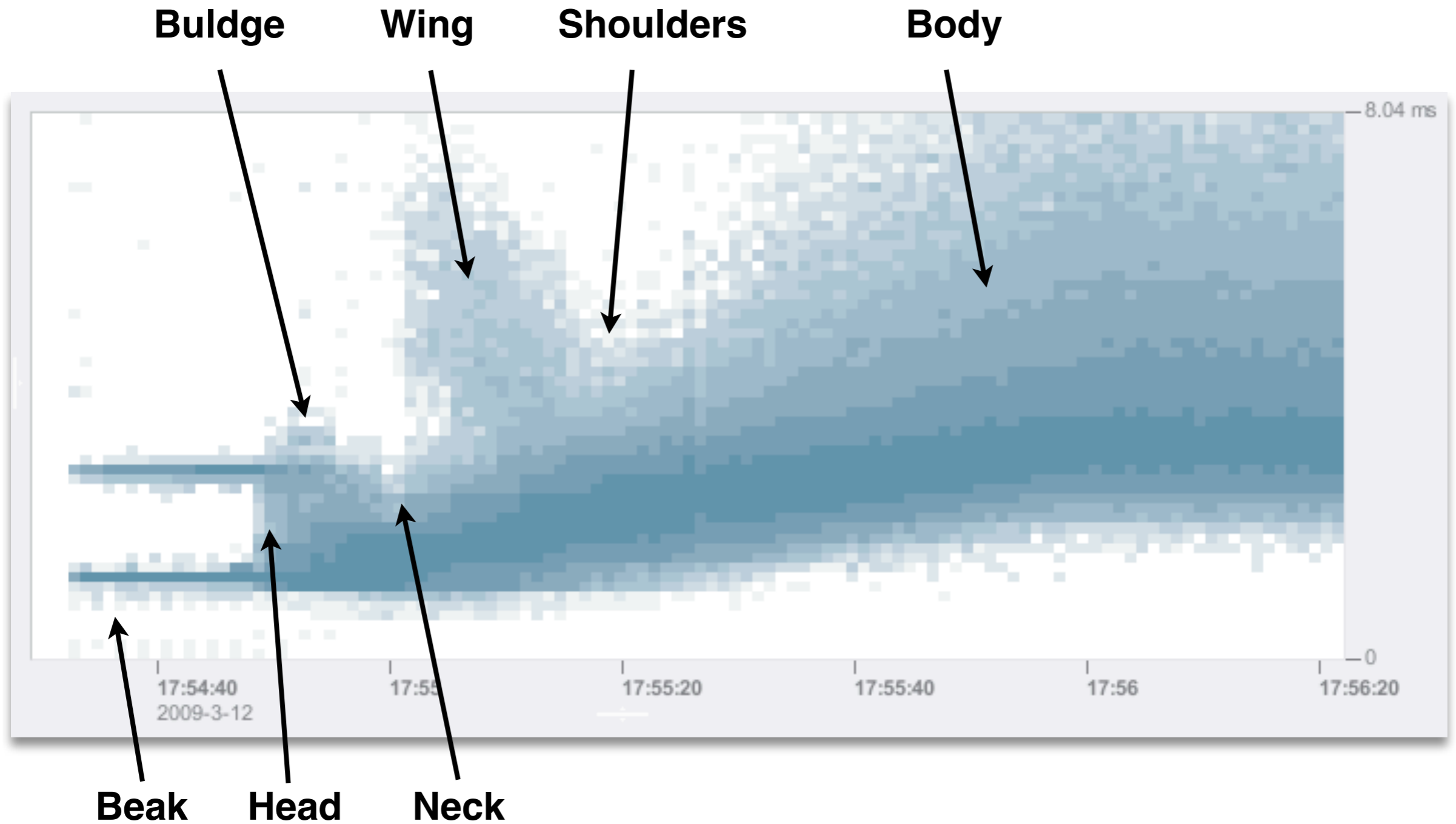
The “Rainbow Pterodactyl”



The “Rainbow Pterodactyl”



The “Rainbow Pterodactyl”



The “Rainbow Pterodactyl”: Analysis



- Hasn't been understood in detail
 - Would never be understood (or even known) without heat maps
- It is repeatable

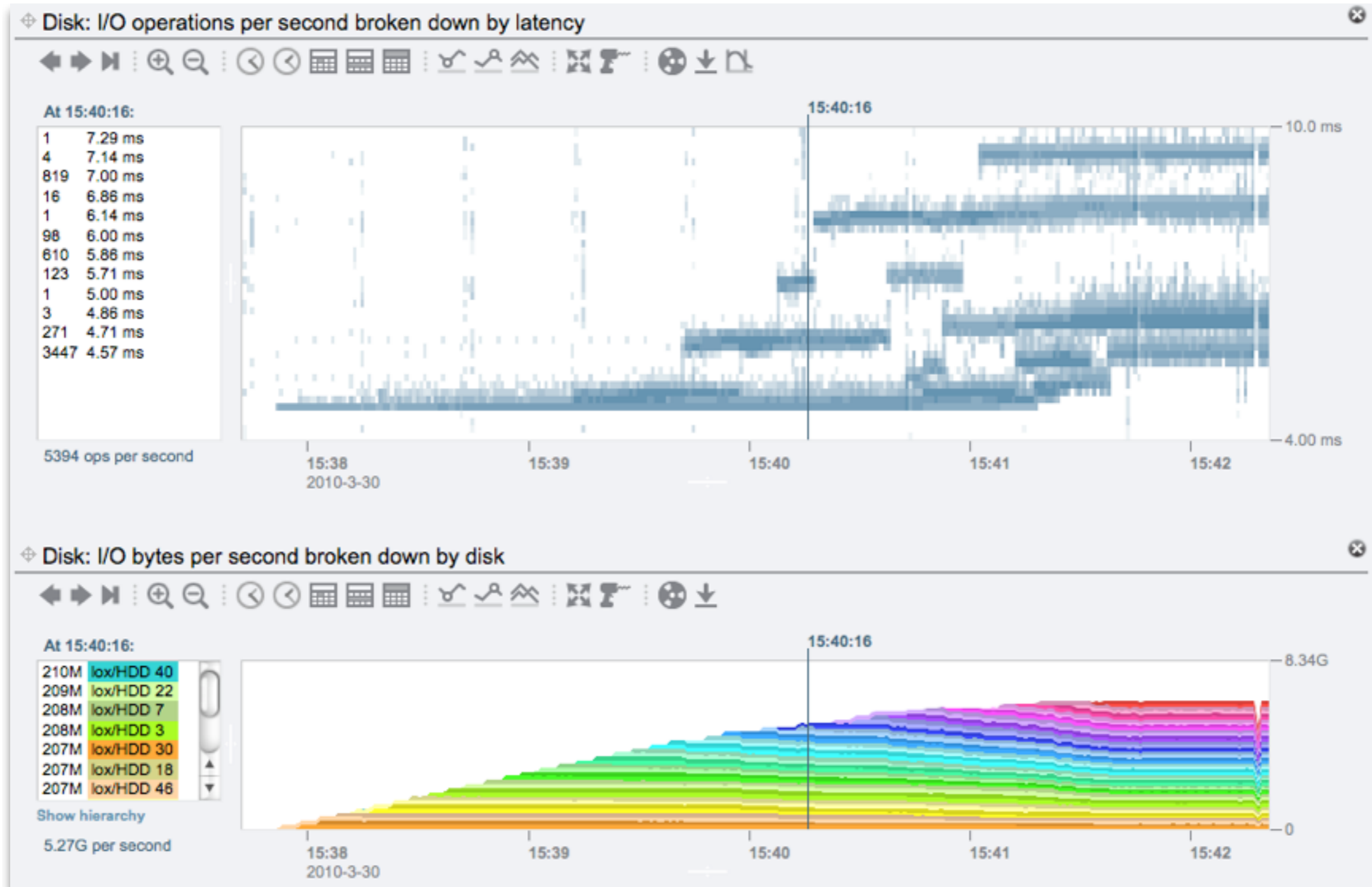
The “Rainbow Pterodactyl”: Theories



- “Beak”: disk cache hit vs disk cache miss -> bimodal
- “Head”: 9th disk, contention on the 2 x4 SAS ports
- “Buldge”: ?
- “Neck”: ?
- “Wing”: contention?
- “Shoulders”: ?
- “Body”: PCI-gen1 bus contention

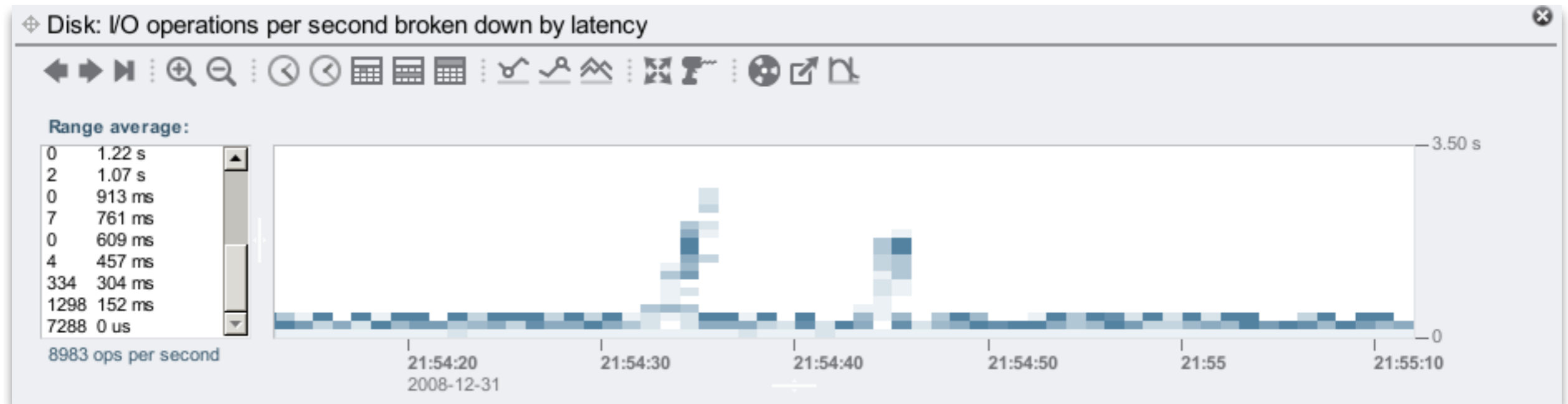
- Same as “Rainbow Pterodactyl”, stepping disks
- Instead of sequential reads, this is repeated 128 Kbyte reads (read -> seek 0 -> read -> ...), to deliberately hit from the disk DRAM cache to improve test throughput

Latency Levels



- ???

Bonus Latency Heat Map



- This time we do know the source of the latency...

癡科技小劇場：硬碟也會鬧情緒

http://chinese.engadget.com/2009/01/05/video-shouting-at-disk-drive-causes-high-latency-low-r

別向硬碟亂叫，會影響它們的心情...

癡科技小劇場：硬碟也會鬧情緒

癡科技小劇場：硬碟也會鬧情緒

由 Flow Yu 於 1 year 之前發表

文章分類: 儲存裝置



這可不是開玩笑來著，而是經過實驗證明的唷！來自 Sun Fishworks 團隊的 Brendan Gregg，或許是在機房裡待久了，於是突發異想跑去對機房中的磁碟陣列大吼大叫一陣，結果就這麼發現了一個真理，那就是：硬碟就跟員工一樣，吼叫並不會讓他們的效率變高，反而還會讓它們心生不爽而降低士氣，私底下就開始搞罷工。

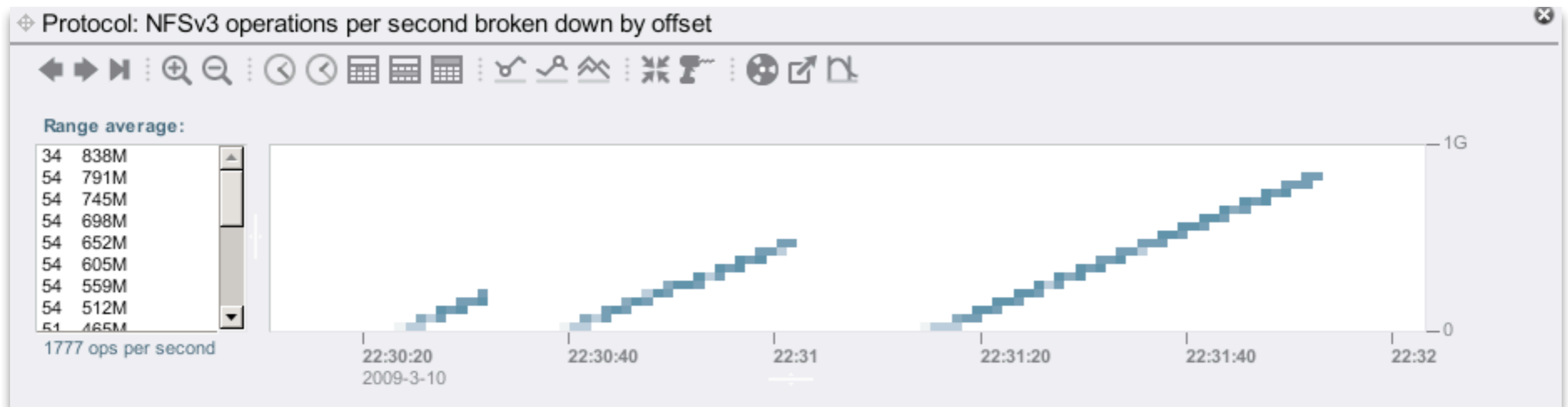
Latency Heat Maps: Summary



- Shows latency distribution over time
- Shows outliers (maximums)
- Indirectly shows average
- Shows patterns
 - allows correlation with other software stack layers

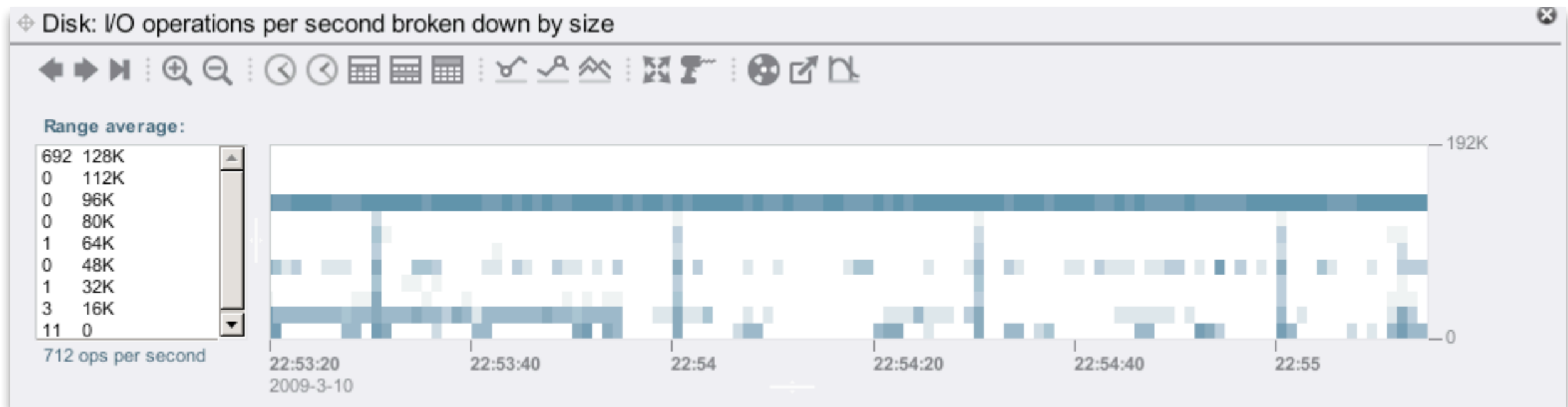
- These all have a dynamic y-axis scale:
 - I/O size
 - I/O offset
- These aren't a primary measure of performance (like latency); they provide secondary information to understand the workload

Heat Map: I/O Offset



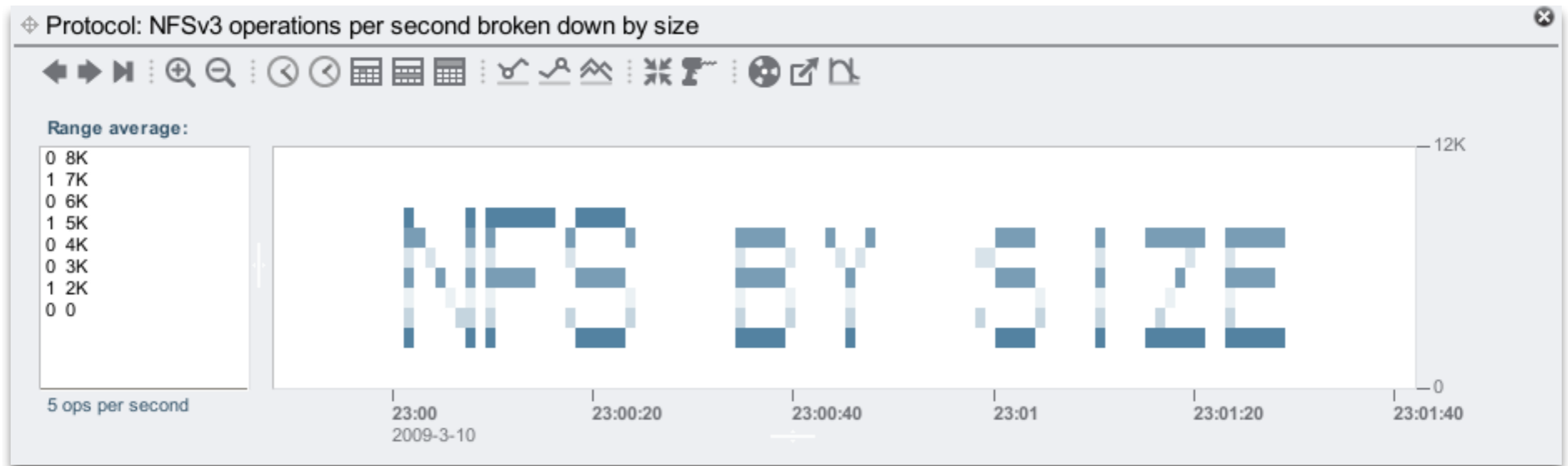
- y-axis: I/O offset (in this case, NFSv3 file location)

Heat Map: I/O Size



- y-axis: I/O size (bytes)

- What can we 'paint' by adjusting the workload?



- How was this done?



- How was this done?



- How was this done?

Current Examples

Utilization



CPU Utilization

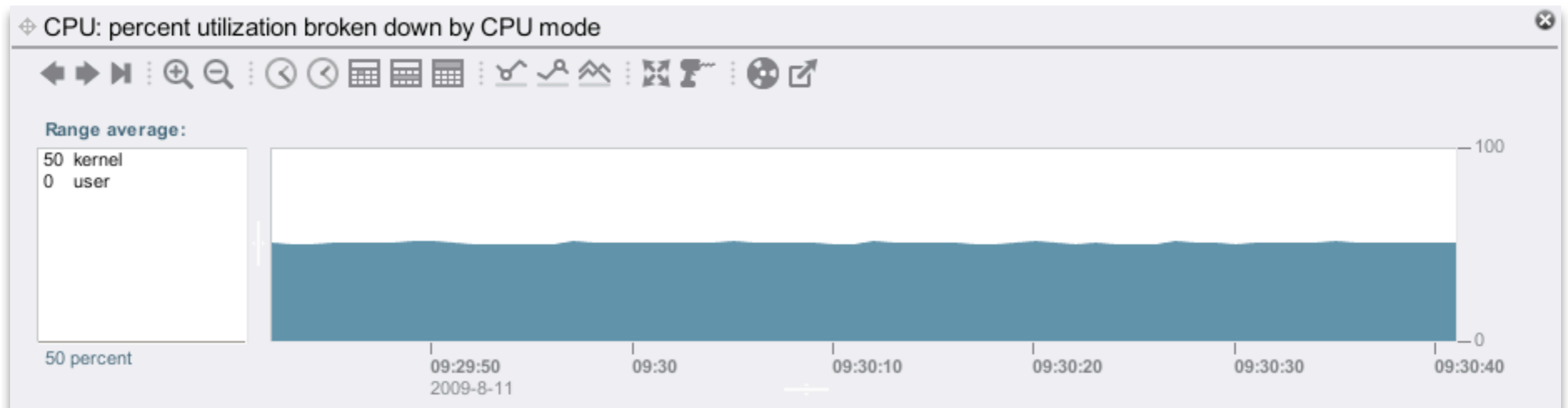


- Commonly used indicator of CPU performance
- eg, vmstat(1M)

```
$ vmstat 1 5
kthr      memory          page        disk        faults      cpu
 r  b  w    swap  free   re  mf pi po fr de sr s0 s1 s2 s3   in   sy   cs us sy id
0  0  0 95125264 28022732 301 1742 1 17 17 0 0 -0 -0 -0  6 5008 21927 3886 4 1 94
0  0  0 91512024 25075924  6  55 0 0  0  0  0  0  0  0  0 4665 18228 4299 10 1 89
0  0  0 91511864 25075796  9  24 0 0  0  0  0  0  0  0  0 3504 12757 3158  8  0 92
0  0  0 91511228 25075164  3 163 0 0  0  0  0  0  0  0  0 4104 15375 3611  9  5 86
0  0  0 91510824 25074940  5  66 0 0  0  0  0  0  0  0  0 4607 19492 4394 10 1 89
```

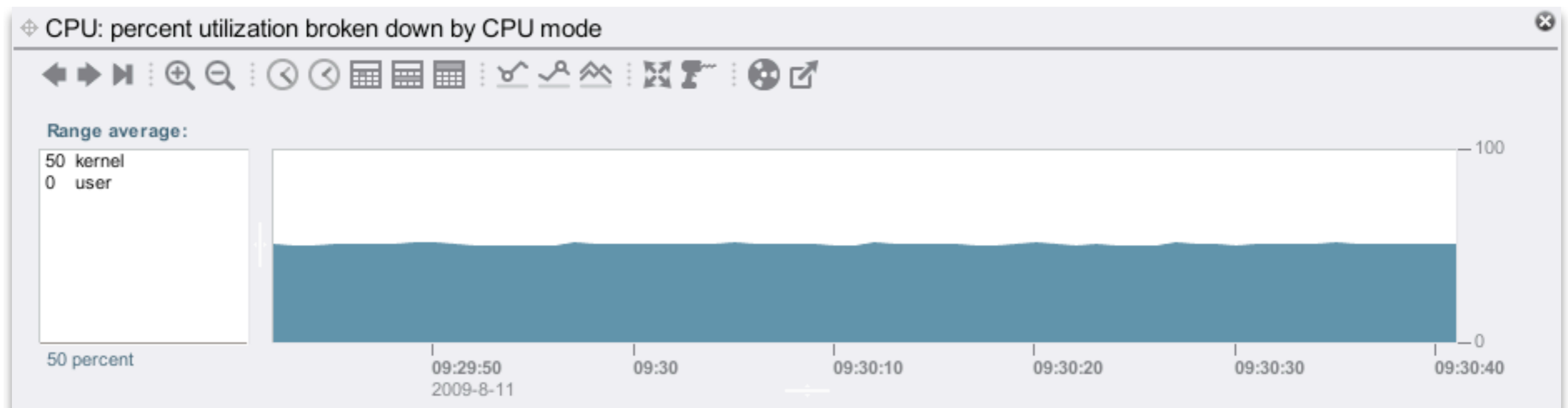
CPU Utilization: Line Graph

- Easy to plot:



CPU Utilization: Line Graph

- Easy to plot:



- Average across all CPUs:
 - Identifies how utilized all CPUs are, indicating remaining headroom - provided sufficient threads to use CPUs

CPU Utilization by CPU



- mpstat(1M) can show utilization by-CPU:

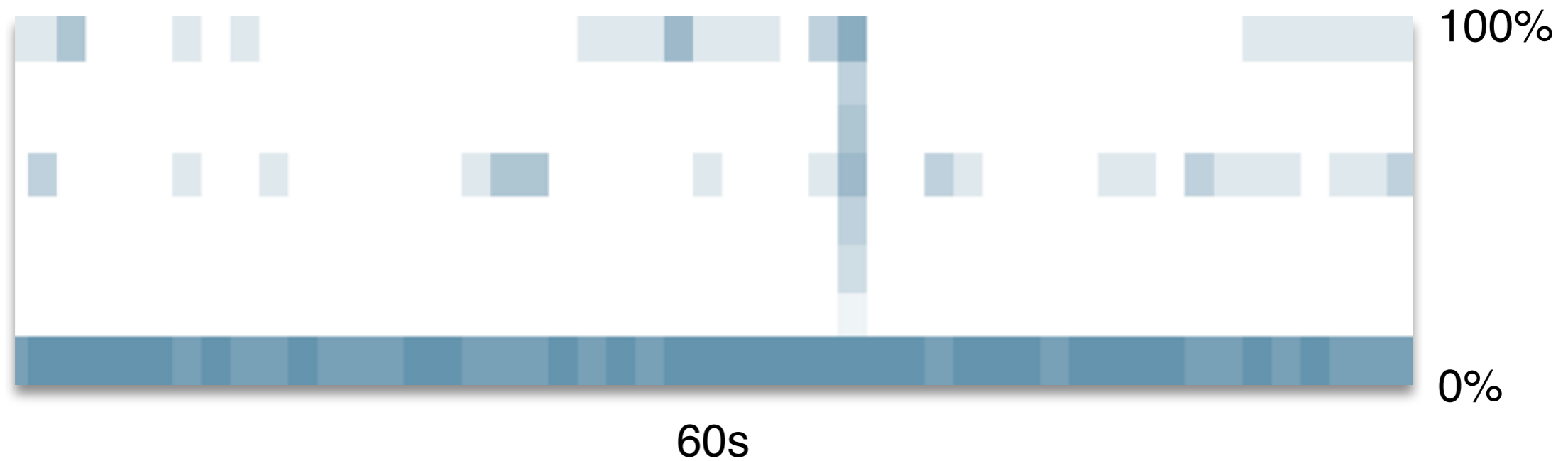
```
$ mpstat 1
[...]
```

CPU	minf	mjf	xcal	intr	ithr	csw	icsw	migr	smtx	srw	syscl	usr	sys	wt	idl
0	0	0	2	313	105	315	0	24	4	0	1331	5	1	0	94
1	0	0	0	65	28	190	0	12	4	0	576	1	1	0	98
2	0	0	0	64	20	152	0	12	1	0	438	0	1	0	99
3	0	0	0	127	74	274	1	21	3	0	537	1	1	0	98
4	0	0	0	32	5	229	0	9	2	0	902	1	1	0	98
5	0	0	0	46	19	138	0	7	3	0	521	1	0	0	99
6	2	0	0	109	32	296	0	8	2	0	1266	4	0	0	96
7	0	0	0	30	8	0	9	0	1	0	0	100	0	0	0
8	0	0	0	169	68	311	0	22	2	0	847	2	1	0	97
9	0	0	30	111	54	274	0	16	4	0	868	2	0	0	98
10	0	0	0	69	29	445	0	13	7	0	2559	7	1	0	92
11	0	0	0	78	36	303	0	7	8	0	1041	2	0	0	98
12	0	0	0	74	34	312	0	10	1	0	1250	7	1	0	92
13	38	0	15	456	285	336	2	10	1	0	1408	5	2	0	93
14	0	0	0	2620	2497	209	0	10	38	0	259	1	3	0	96
15	0	0	0	20	8	10	0	4	2	0	2	0	0	0	100

- can identify a single hot CPU (thread)
 - and un-balanced configurations

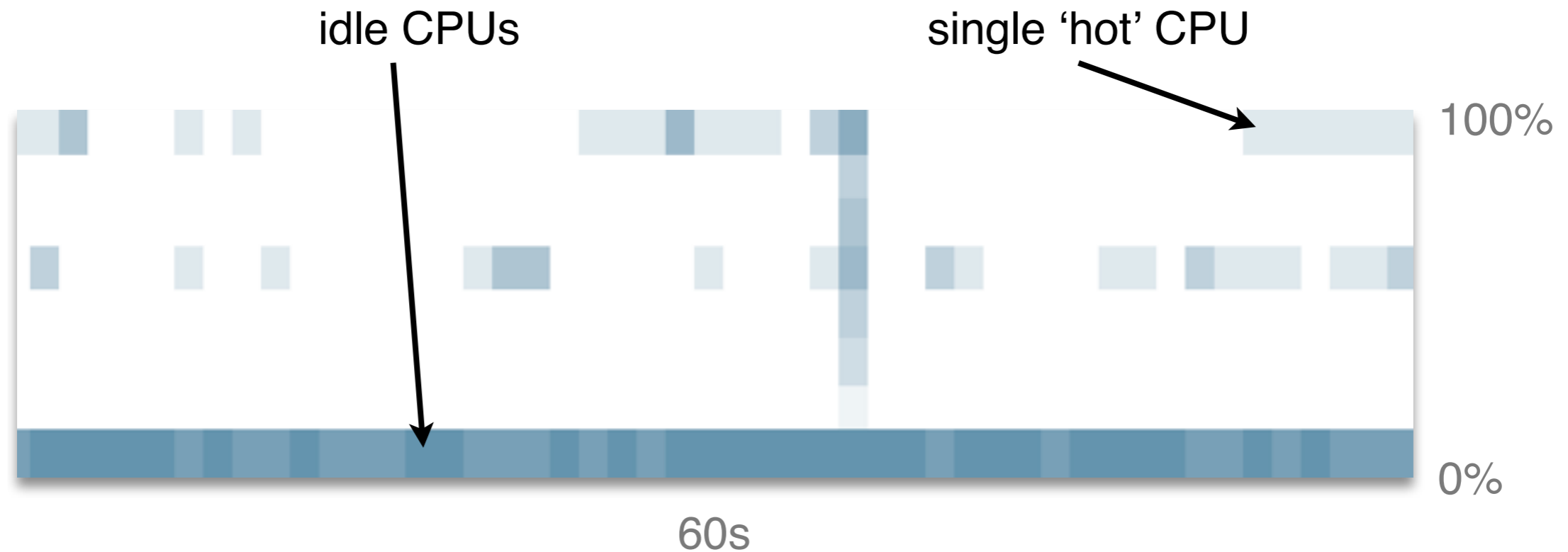
- Monitor overall utilization for capacity planning
- Also valuable to monitor individual CPUs
 - can identify un-balanced configurations
 - such as a single hot CPU (thread)
- The virtual CPUs on a single host can now reach the 100s
 - its own dimension
 - how can we display this 3rd dimension?

Heat Map: CPU Utilization



- x-axis: time
- y-axis: percent utilization
- z-axis (color saturation): # of CPUs in that time/utilization range

Heat Map: CPU Utilization



- Single 'hot' CPUs are a common problem due to application scalability issues (single threaded)
- This makes identification easy, without reading pages of mpstat (1M) output

- Ditto for disks
- Disk Utilization heat map can identify:
 - overall utilization
 - unbalanced configurations
 - single hot disks (versus all disks busy)
- Ideally, the disk utilization heat map is tight (y-axis) and below 70%, indicating a well balanced config with headroom
 - which can't be visualized with line graphs

- Are typically used to visualize performance, be it IOPS or utilization
- Show patterns over time more clearly than text (higher resolution)
- But graphical environments can do much more
 - As shown by the heat maps (to start with); which convey details line graphs cannot
- Ask: what “value add” does the GUI bring to the data?

- Can exist for any resource with multiple components:
 - CPUs
 - Disks
 - Network interfaces
 - I/O busses
 - ...
- Quickly identifies single hot component versus all components
- Best suited for physical hardware resources
 - difficult to express 'utilization' for a software resource

Future Opportunities



So far analysis has been for a single server

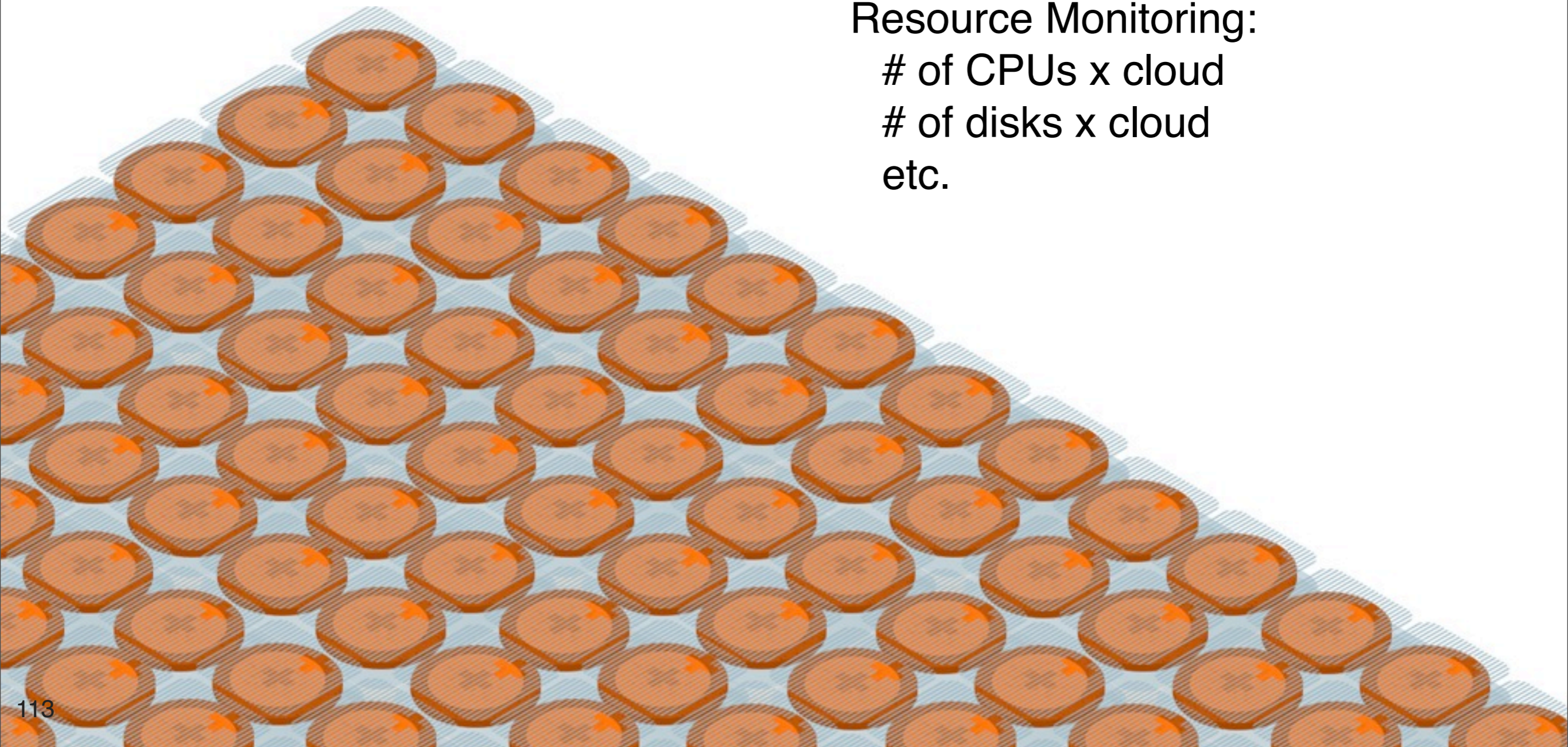
What about the cloud?



From one to thousands of servers

Workload Analysis:
latency I/O x cloud

Resource Monitoring:
of CPUs x cloud
of disks x cloud
etc.



- Heat Maps are promising for cloud computing observability:
 - additional dimension accommodates the scale of the cloud
- Find outliers regardless of node
 - cloud-wide latency heat map just has more I/O
- Examine how applications are load balanced across nodes
 - similar to CPU and disk utilization heat maps
- mpstat and iostat's output are already getting too long
 - multiply by 1000x for the number of possible hosts in a large cloud application

- Include:
 - Latency heat map across entire cloud
 - Latency heat maps for cloud application components
 - CPU utilization by cloud node
 - CPU utilization by CPU
 - Thread/process utilization across entire cloud
 - Network interface utilization by cloud node
 - Network interface utilization by port
 - lots, lots more

- Latency at different layers:
 - Apache
 - PHP/Ruby/...
 - MySQL
 - DNS
 - Disk I/O
 - CPU dispatcher queue latency
 - and pattern match to quickly identify and locate latency

Latency Example: MySQL



- Query latency (DTrace):

query time (ns)	value	----- Distribution -----	count
	1024		0
	2048		2
	4096	@	99
	8192		20
	16384	@	114
	32768	@	105
	65536	@	123
	131072	@@@@@@@@@@@@@@@@	1726
	262144	@@@@@@@@@@@@@@	1515
	524288	@@@	601
	1048576	@@	282
	2097152	@	114
	4194304		61
	8388608	@@@@	660
	16777216		67
	33554432		12
	67108864		7
	134217728		4
	268435456		5
	536870912		0

Latency Example: MySQL



- Query latency (DTrace):

query time (ns)	value	----- Distribution -----	count
	1024		0
	2048		2
	4096	@	99
	8192		20
	16384	@	114
	32768	@	105
	65536	@	123
	131072	@@@@@@@@@@@@@@@@	1726
	262144	@@@@@@@@@@@@@@@@	1515
	524288	@@@	601
	1048576	@@	282
	2097152	@	114
	4194304		61
	8388608	@@@@@	660
	16777216		67
	33554432		12
	67108864		7
	134217728		4
	268435456		5
	536870912		0

**What is this?
(8-16 ms latency)**

Latency Example: MySQL



query time (ns)	value	----- Distribution -----	count
	1024		0
	2048		2
	4096	@	99
	8192		20
	16384	@	114
	32768	@	105
	65536	@	123
	131072	@@@@@@@@@@@@@@@@	1726
	262144	@@@@@@@@@@@@@@@@	1515
	524288	@@@	601
	1048576	@@	282
	2097152	@	114
	4194304		61
	8388608	@@@@	660
	16777216		67
	33554432		12
	67108864		7
	134217728		4
	268435456		5
	536870912		0

oh...

innodb srv sleep (ns)	value	----- Distribution -----	count
	4194304		0
	8388608	@@	841
	16777216		0

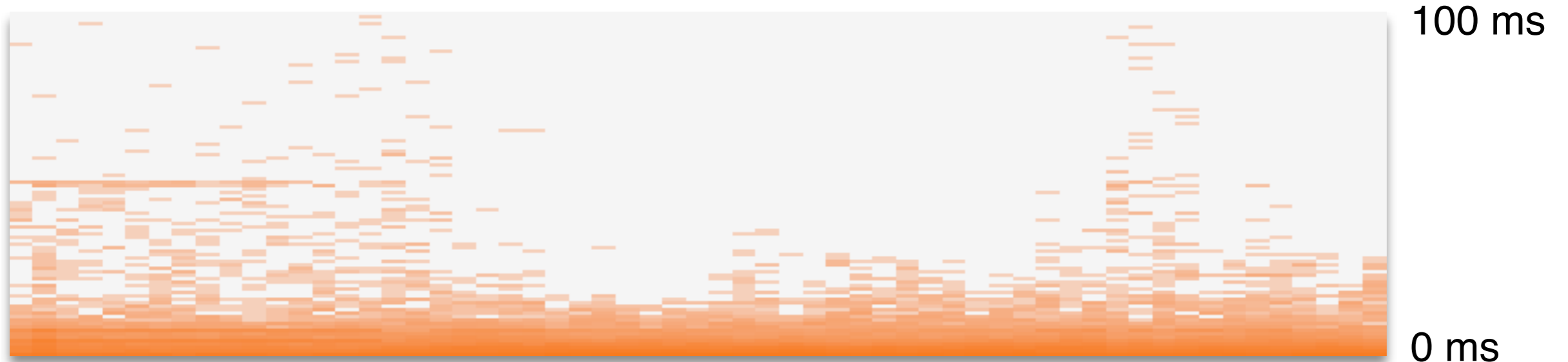
- Spike of MySQL query latency: 8 - 16 ms
- innodb thread concurrency back-off sleep latency: 8 - 16 ms
- Both have a similar magnitude (see “count” column)
- Add the dimension of time as a heat map, for more characteristics to compare
- ... quickly compare heat maps from different components of the cloud to pattern match and locate latency

- Identify latency outliers, distributions, patterns
- Can add more functionality to identify these by:
 - cloud node
 - application, cloud-wide
 - I/O type (eg, query type)
- Targeted observability (DTrace) can be used to fetch this
- Or, we could collect it for everything
 - ... do we need a 4th dimension?

4th Dimension!

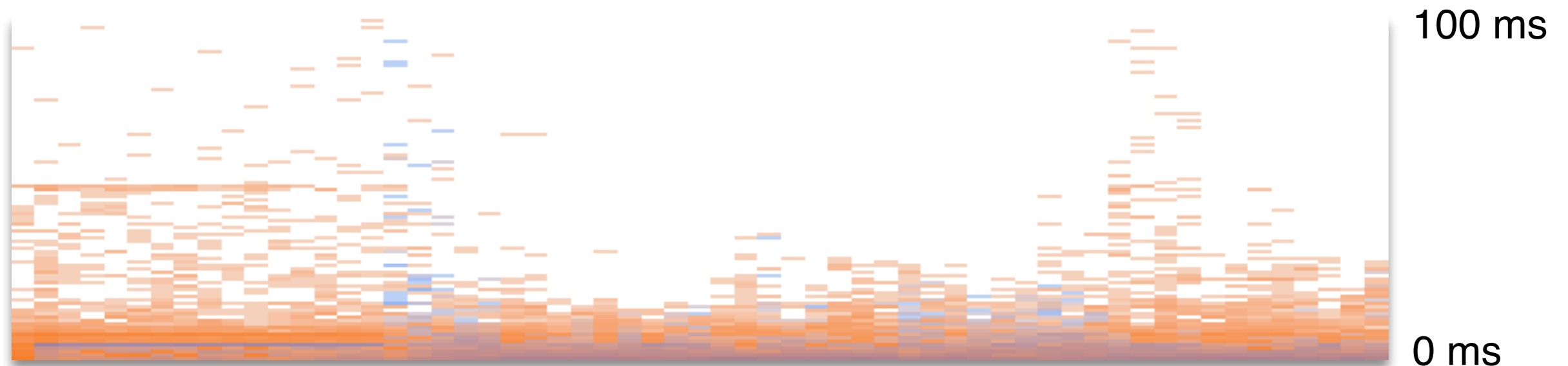
- Bryan Cantrill @Joyent coded this 11 hours ago
 - assuming it's now about 10:30am during this talk
 - ... and I added these slides about 7 hours ago

4th Dimension Example: Thread Runtime



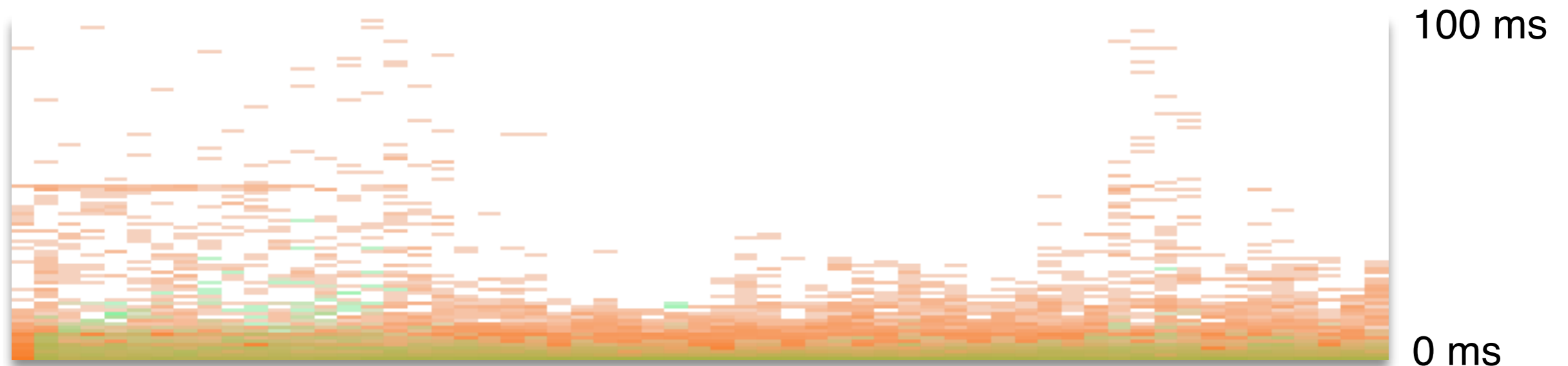
- x-axis: time
- y-axis: thread runtime
- z-axis (color saturation): count at that time/runtime range

4th Dimension Example: Thread Runtime



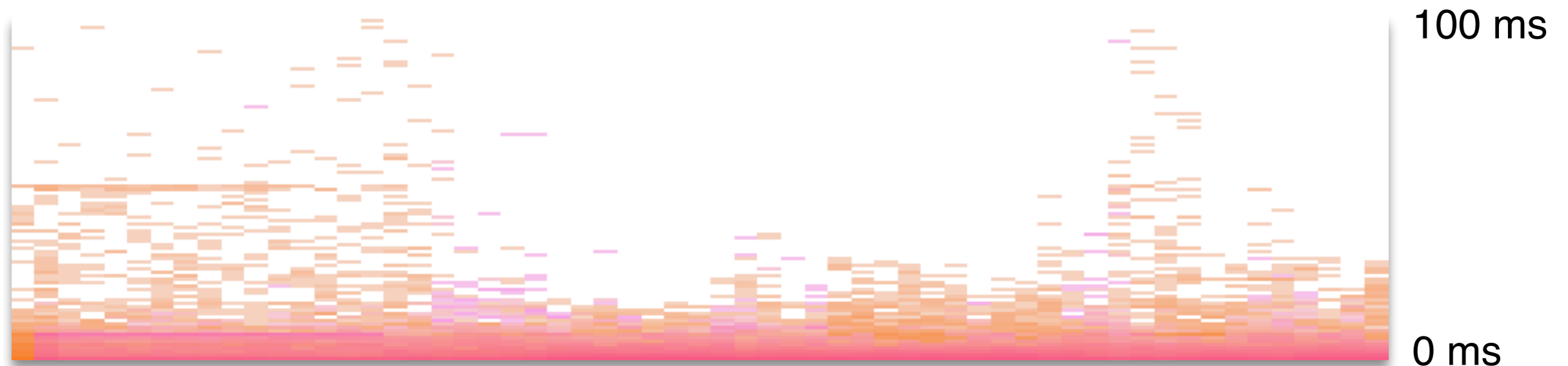
- x-axis: time
- y-axis: thread runtime
- z-axis (color saturation): count at that time/runtime range
- omega-axis (color hue): application
- blue == “coreaudiod”

4th Dimension Example: Thread Runtime



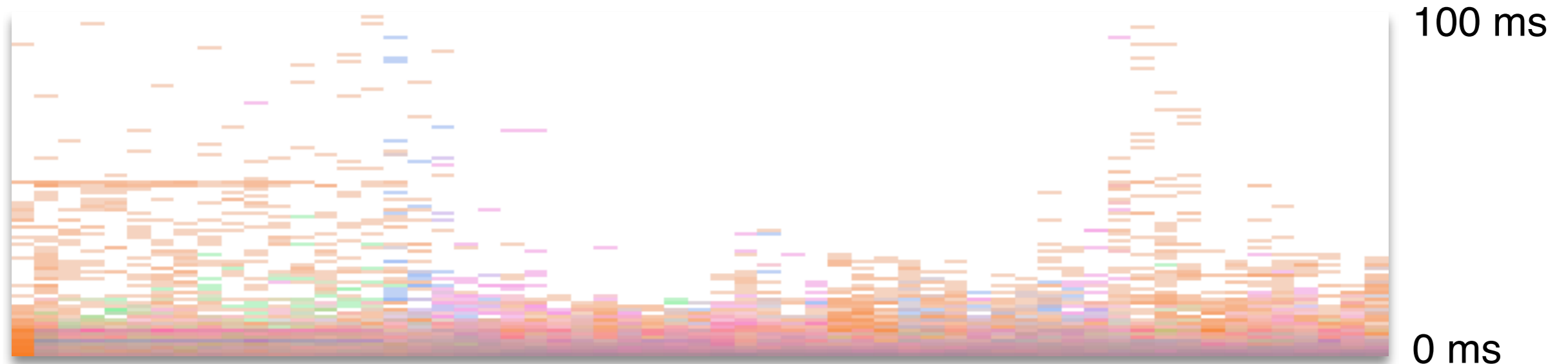
- x-axis: time
- y-axis: thread runtime
- z-axis (color saturation): count at that time/runtime range
- omega-axis (color hue): application
- green == "iChat"

4th Dimension Example: Thread Runtime



- x-axis: time
- y-axis: thread runtime
- z-axis (color saturation): count at that time/runtime range
- omega-axis (color hue): application
- violet == "Chrome"

4th Dimension Example: Thread Runtime



- x-axis: time
- y-axis: thread runtime
- z-axis (color saturation): count at that time/runtime range
- omega-axis (color hue): application
- All colors

- While the data supports the 4th dimension, visualizing this properly may become difficult (we are eager to find out)
 - The image itself is still only 2 dimensional
- May be best used to view a limited set, to limit the number of different hues; uses can include:
 - Highlighting different cloud application types: DB, web server, etc.
 - Highlighting one from many components: single node, CPU, disk, etc.
- Limiting the set also helps storage of data

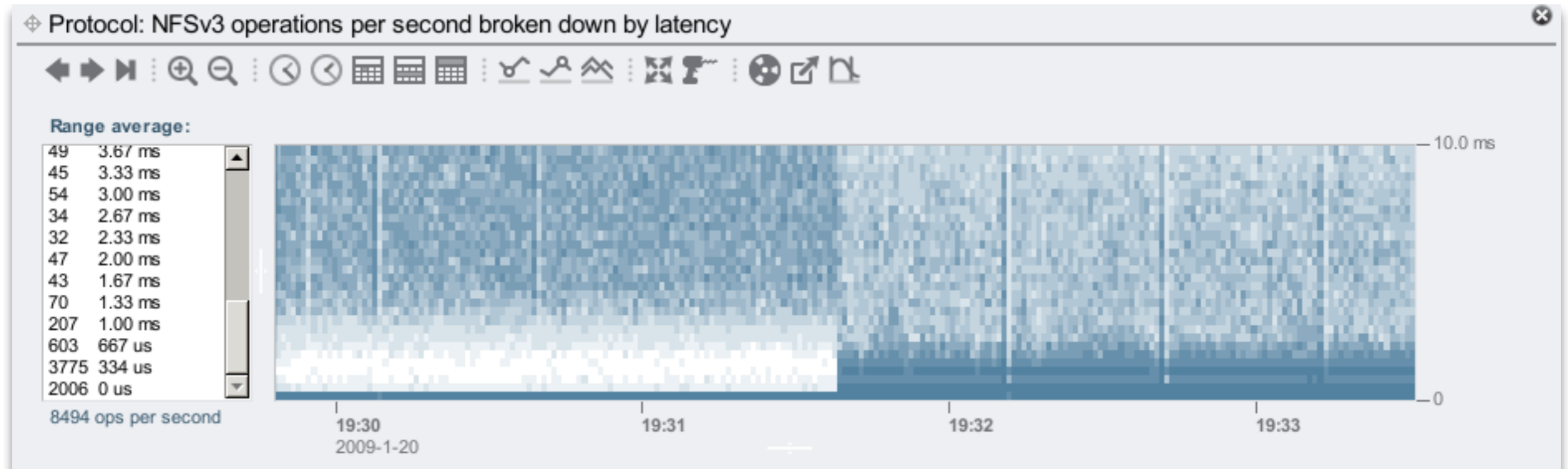
- We plan much more new stuff
 - We are building a team of engineers to work on it; including Bryan Cantrill, Dave Pacheco, and mysqlf
 - Dave and I have only been at Joyent for 2 1/2 weeks

- Visualizations such as heat maps could also be applied to:
- Security, with pattern matching for:
 - robot identification based on think-time latency analysis
 - inter-keystroke-latency analysis
 - brute force username latency attacks?
- System Administration
 - monitoring quota usage by user, filesystem, disk
- Other multi-dimensional datasets

- Consider performance metrics before plotting
 - Why latency is good
 - ... and IOPS can be bad
- See the value of visualizations
 - Why heat maps are needed
 - ... and line graphs can be bad
- Remember key examples
 - I/O latency, as a heat map
 - CPU utilization by CPU, as a heat map

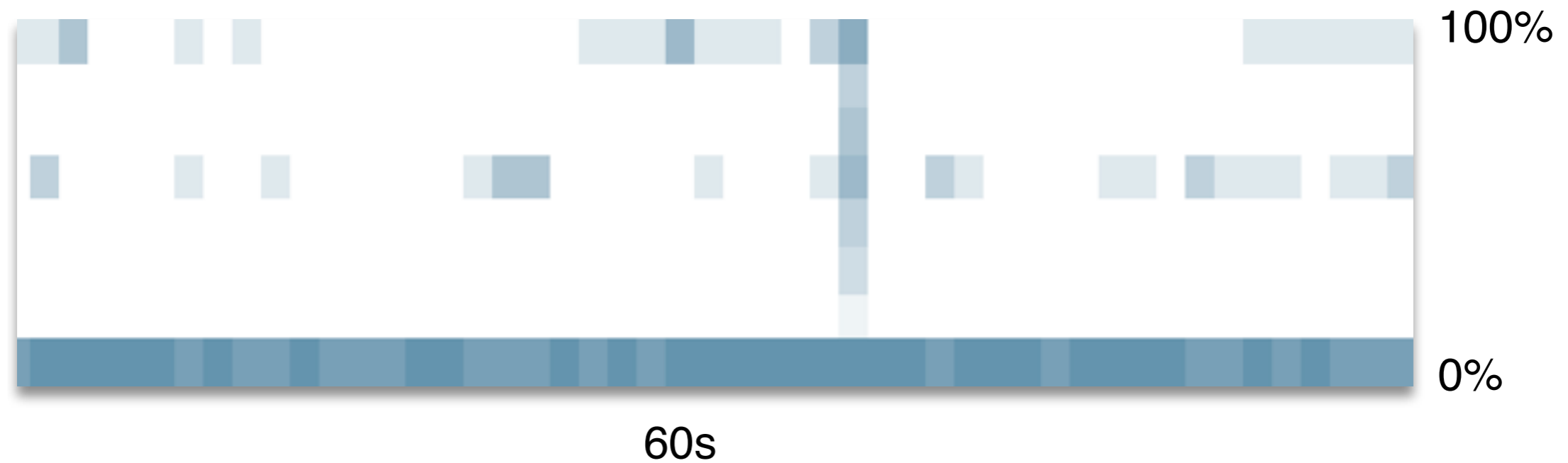


Heat Map: I/O Latency



- Latency matters
 - synchronous latency has a direct impact on performance
- Heat map shows
 - outliers, balance, cache layers, patterns

Heat Map: CPU Utilization



- Identify single threaded issues
 - single CPU hitting 100%
- Heat map shows
 - fully utilized components, balance, overall headroom, patterns

- For Reference:
- DTraceTazTool
 - 2006; based on TazTool by Richard McDougall 1995. Open source, unsupported, and probably no longer works (sorry).
- Analytics
 - 2008; Oracle Sun ZFS Storage Appliance
- “new stuff” (not named yet)
 - 2010; Joyent; Bryan Cantrill, Dave Pacheco, Brendan Gregg

- Thank you!

- How to find me on the web:
 - <http://dtrace.org/blogs/brendan>
 - <http://blogs.sun.com/brendan> <-- is my old blog
 - twitter @brendangregg